

# Your Time Is Now



# SR Traffic Engineering

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BRKRST-3122

Strictly Confidential

# Industry at large backs up SR



**Strong customer  
adoption**

WEB, SP, Enterprise



**Standardization**  
IETF



**De-Facto SDN  
Architecture**



**Multi-vendor  
Consensus**  
Interop testings

# Stay Up-To-Date

[amzn.com/B01I58LSUO](https://amzn.com/B01I58LSUO)



[segment-routing.net](https://segment-routing.net)



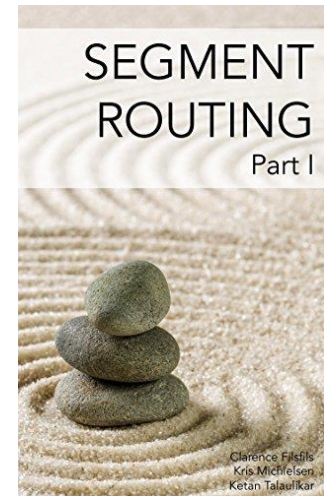
[linkedin.com/groups/8266623](https://linkedin.com/groups/8266623)



[twitter.com/SegmentRouting](https://twitter.com/SegmentRouting)



[facebook.com/SegmentRouting/](https://facebook.com/SegmentRouting/)



# IETF key document for SR-TE

Network Working Group  
Internet-Draft  
Intended status: Standards Track  
Expires: August 22, 2017

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February 18, 2017

Segment Routing Policy for Traffic Engineering  
draft-filsfils-spring-segment-routing-policy-00.txt

# SR Traffic Engineering

SR-TE

# RSVP-TE

- Little deployment and many issues
- Not scalable
  - Core states in  $k \cdot n^2$
  - No inter-domain
- Complex configuration
  - Tunnel interfaces
- Complex steering
  - PBR, autoroute

# SR TE

- Simple, Automated and Scalable
  - No core state: state in the packet header
  - No tunnel interface: “SR Policy”
  - No headend a-priori configuration: on-demand policy instantiation
  - No headend a-priori steering: on-demand steering
- Multi-Domain
  - XTC for compute
  - BSID for scale
- Lots of Functionality
  - Designed with lead operators along their use-cases



# SR Policy

# SR Policy

```
segment-routing
traffic-eng
policy FOO
  end-point ipv4 1.1.1.4 color 20
  binding-sid mpls 1000
  path
    preference 100
    explicit SIDLIST1
  preference 200
    dynamic mpls
    metric
      type latency
    affinity
      exclude-any red
  explicit-path name SIDLIST1
    index 10 mpls label 16002
    index 20 mpls label 30203
    index 30 mpls label 16004
```

SR policy (1.1.1.4, 20)

Path received via BGP signaling  
preference 300  
binding-sid mpls 1000  
weight 1, SID list <16002, 16005>  
weight 2, SID list <16004, 16008>

Path received via PCEP signaling  
preference 400  
binding-sid mpls 1000  
SID list <16002, 16005>

Path received via NETCONF signaling  
preference 500  
binding-sid mpls 1000  
SID list <16002, 16005>

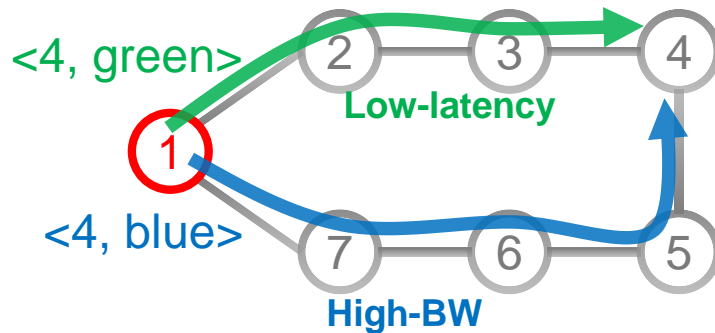
## FIB @ headend

Incoming label: 1000

Action: pop and push <16002, 30203, 14004>

# SR Policy

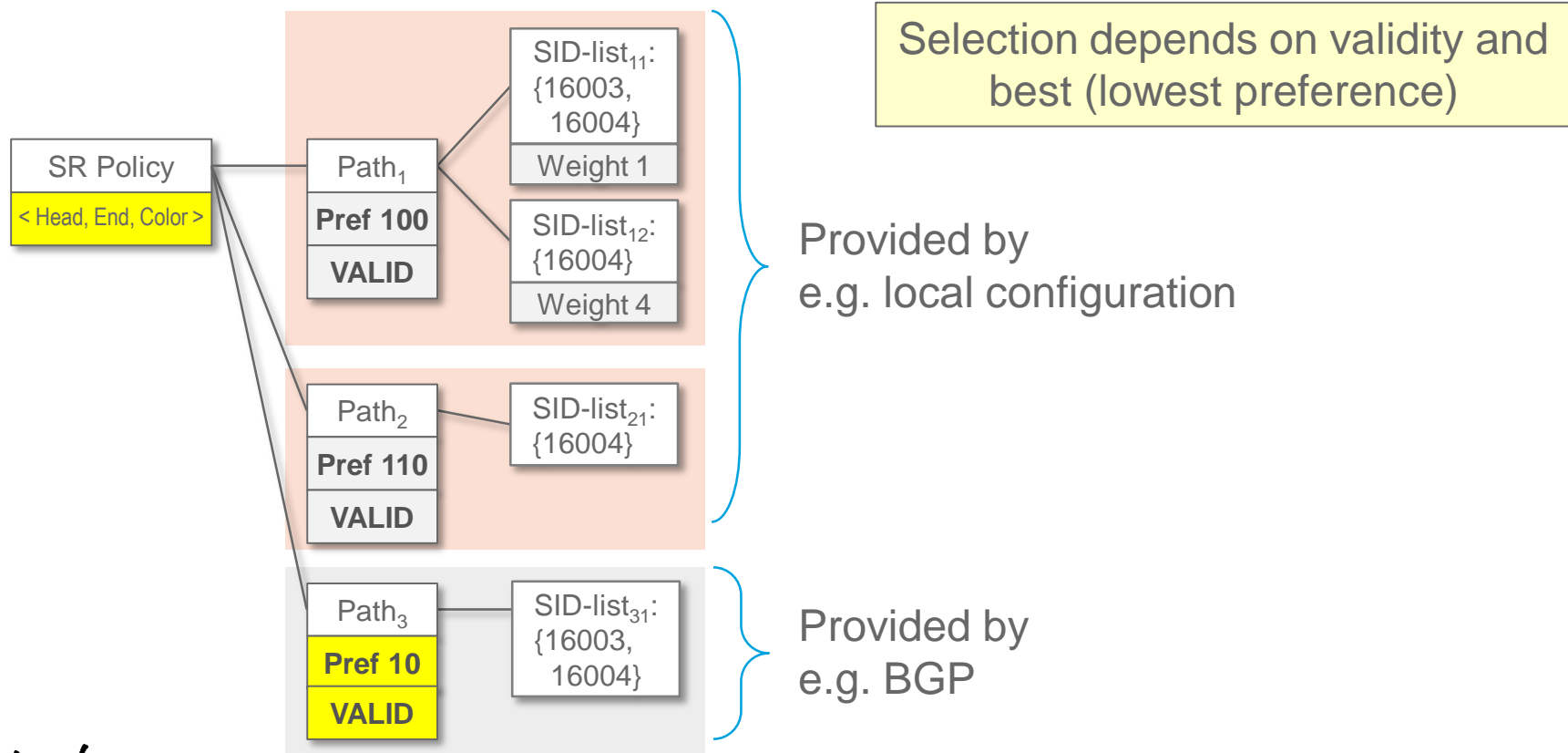
- An SR Policy is identified through the following tuple:
  - The **head-end** where the policy is instantiated/implemented
  - The **endpoint** (i.e.: the destination of the policy)
  - The **color** (an arbitrary numerical value)
- At a given head-end, an SR Policy is fully identified by the **<color, endpoint>** tuple
- An endpoint can be specified as an IPv4 or IPv6 address



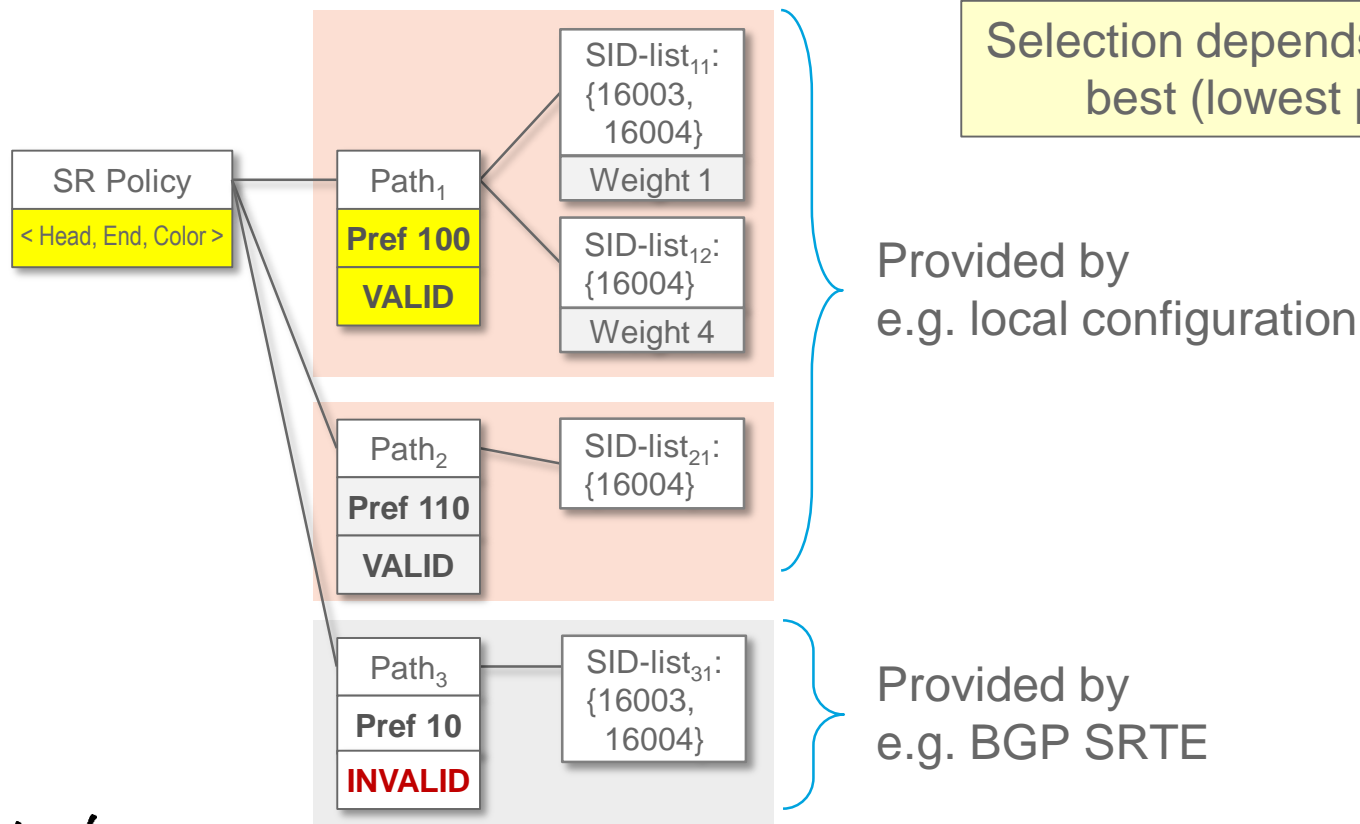
# SRTE DB

- A headend can learn an attached domain topology via its IGP or a BGP-LS session
- A headend can learn a non-attached domain topology via a BGP-LS session
- A headend collects all these topologies in the SR-TE database (SRTE-DB).
- The SRTE-DB is multi-domain capable

# Path's source does not influence selection



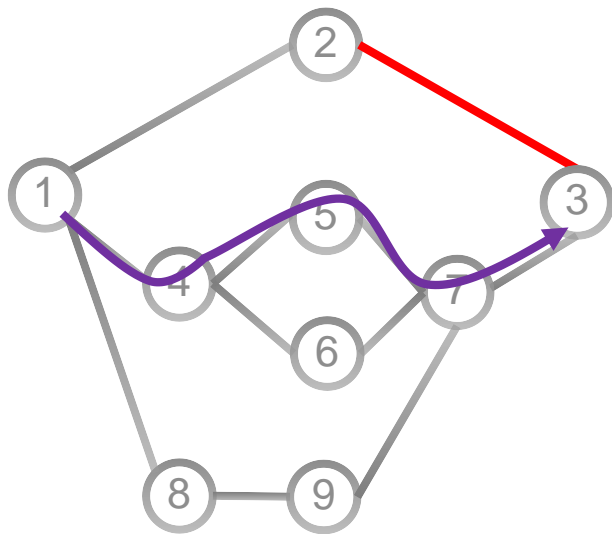
# Path's source does not influence selection



# Dynamic Path

# Headend Computation

# Prefer SR-native Algorithm

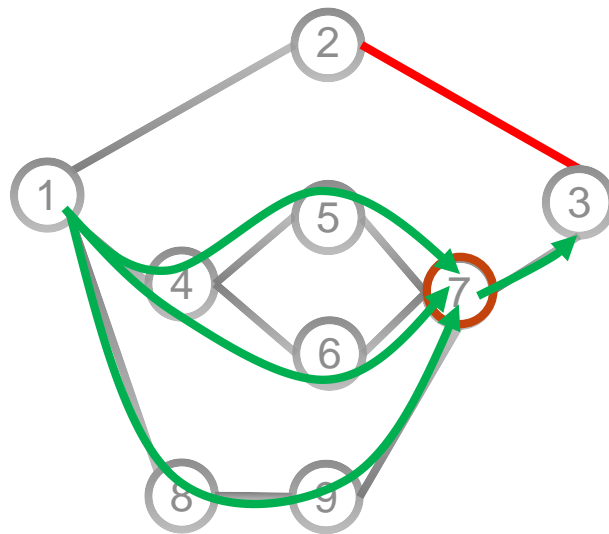


Classic Circuit Algo is not optimum!

SID List: {4, 5, 7, 3}

Poor ECMP, big SR list

ATM optimized



SR-native is optimum

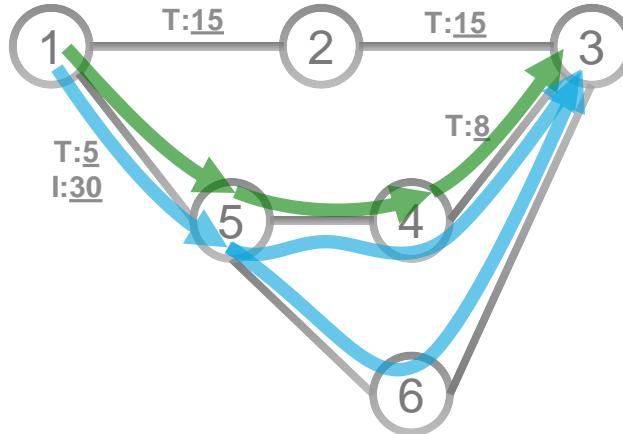
Shortest SID list with Max ECMP

SID List: {7, 3}

IP-optimized



# Min-Metric with Margin



Default IGP link metric: I:10  
Default TE link metric: T:10

Min-Metric(1 to 3, TE)  
= SID-list <16005, 16004, 16003>  
Cumulated TE metric = 23

```
segment-routing
traffic-eng
policy F00
end-point ipv4 1.1.1.3 color 20
binding-sid mpls 1000
path
  preference 50
  dynamic mpls
  metric
    type te
    margin 5
    sid-limit 6
```

Min-Metric(1 to 3, TE, m=5, s<=6)  
= SID-list <16005, 16003>  
Max Cumulated TE metric = 25 < 23+ 5

```

segment-routing
  traffic-eng
    policy FOO
      end-point 1.2.3.4 color 10
      path
        preference 100
        affinity
          include-any RED
          exclude-any BLACK
        address
          include PFXSET1
          exclude PFXSET2
        srlg
          include 123
          exclude 654
        admin-tag
          include 1111
          exclude 3333
      !
      dynamic mpls
        metric
          type te
          limit 200
          sid-limit 5
          sid-list-limit 1
          association group 1
type node

```

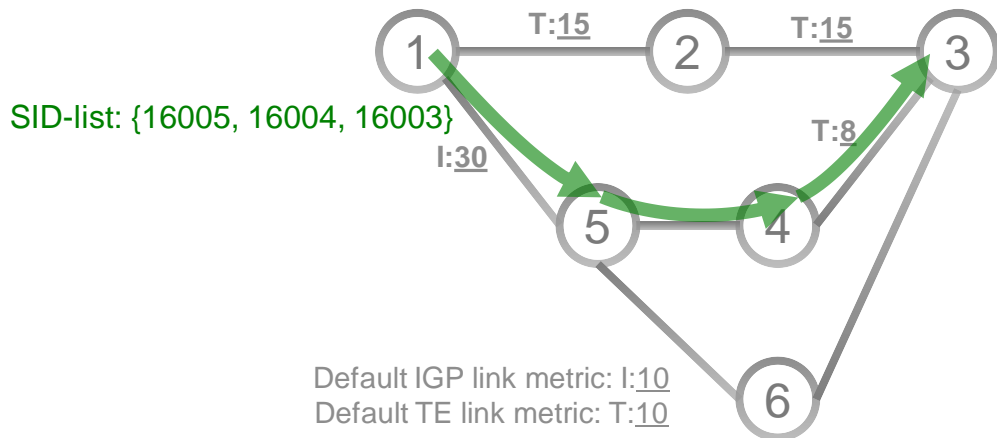
```

segment-routing
  traffic-eng
    policy FOO
      end-point ipv4 1.1.1.3 color 20
      binding-sid mpls 1000
      path
        preference 50
        dynamic mpls
          metric
            type te
            margin 10
            sid-limit 6
      constraints

```

Headend computes a  
SID list respecting  
these constraints

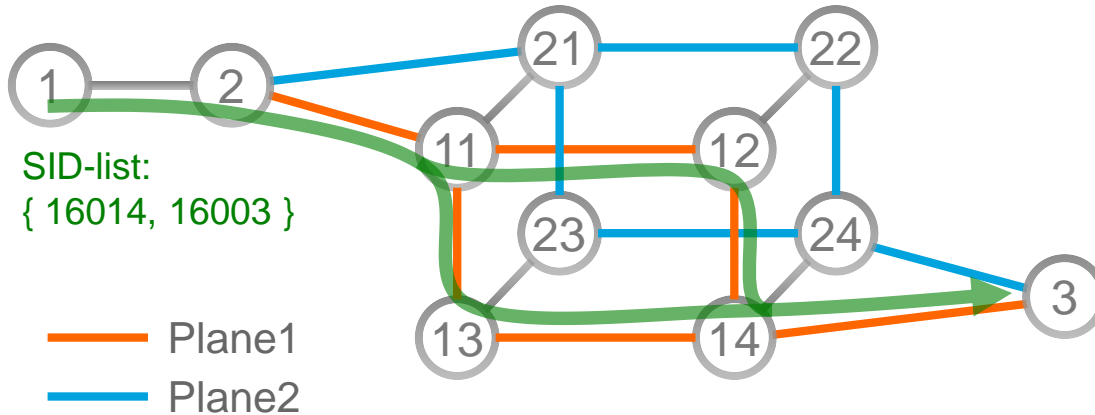
# Low-Latency



```
segment-routing
traffic-eng
policy F00
  end-point ipv4 1.1.1.3 color 20
  binding-sid mpls 1000
  path
    preference 50
    dynamic mpls
    metric
      type te
```

- Min-metric on TE metric where propagation latency is encoded in TE metric
  - same with margin and Max-SID
  - same with latency metric automatically measured by a node for its attached links and distributed in the IGP

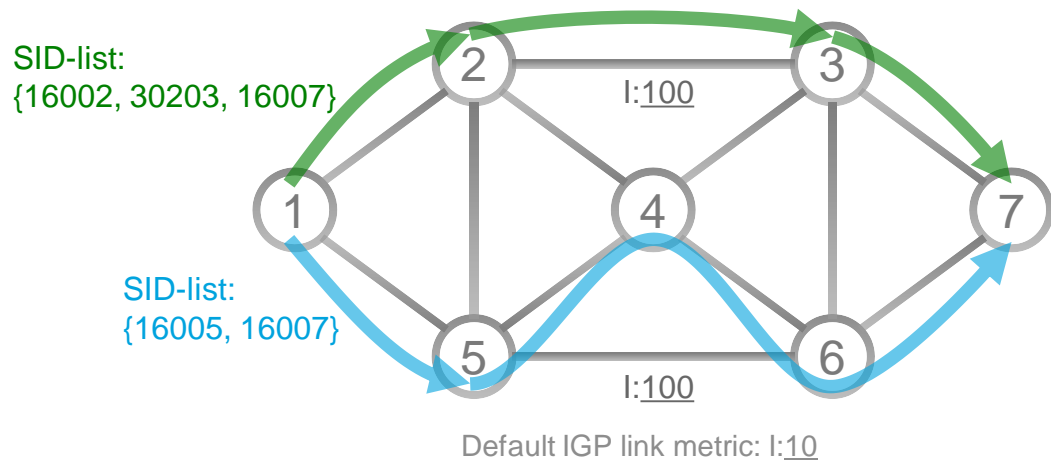
# Plane Affinity



```
segment-routing
traffic-eng
policy F00
  end-point ipv4 1.1.1.3 color 30
  binding-sid mpls 1001
  path
    preference 50
    affinity
      exclude-any Plane2
    dynamic mpls
    metric
      type igp
```

- Min-Metric on IGP metric with exclusion of a TE-affinity “Plane2”
  - all the links part of plane 2 are set with TE-affinity “Plane2”

## Service Disjointness from same headend



```
segment-routing
traffic-eng
policy POLICY1
  end-point ipv4 1.1.1.7 color 100
  path
  preference 50
  dynamic mpls
  metric
  type igp
  association group 1 type node

policy POLICY2
  end-point ipv4 1.1.1.7 color 200
  path
  preference 50
  dynamic mpls
  metric
  type igp
  association group 1 type node
```

- The headend computes two disjoint paths

# On-demand SR Policy

Intra-Domain

# On-Demand SR Policy

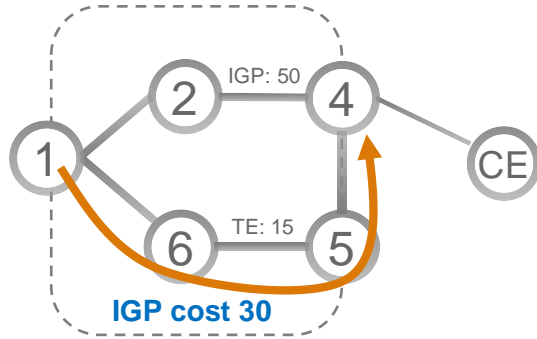
- A service head-end **automatically instantiates** an SR Policy to a BGP nhop when required (on-demand), **automatically steering** the BGP traffic into this SR Policy
- Color community is used as SLA indicator
- Reminder: an SR policy is defined (endpoint, color)

BGP  
Next-hop

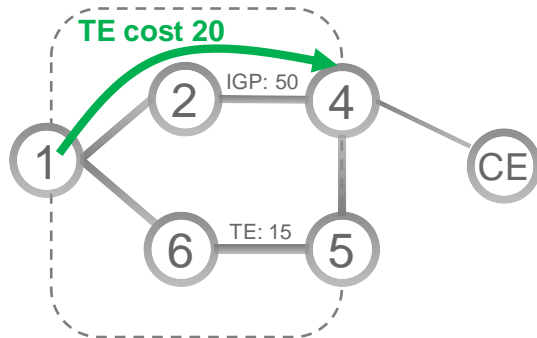


BGP Color  
Community

# Different VPNs need different underlay SLA



Basic VPN should  
use lowest cost  
underlay path



Premium VPN  
should use lowest  
latency path

Objective:  
operationalize  
this service for  
simplicity, scale  
and  
performance



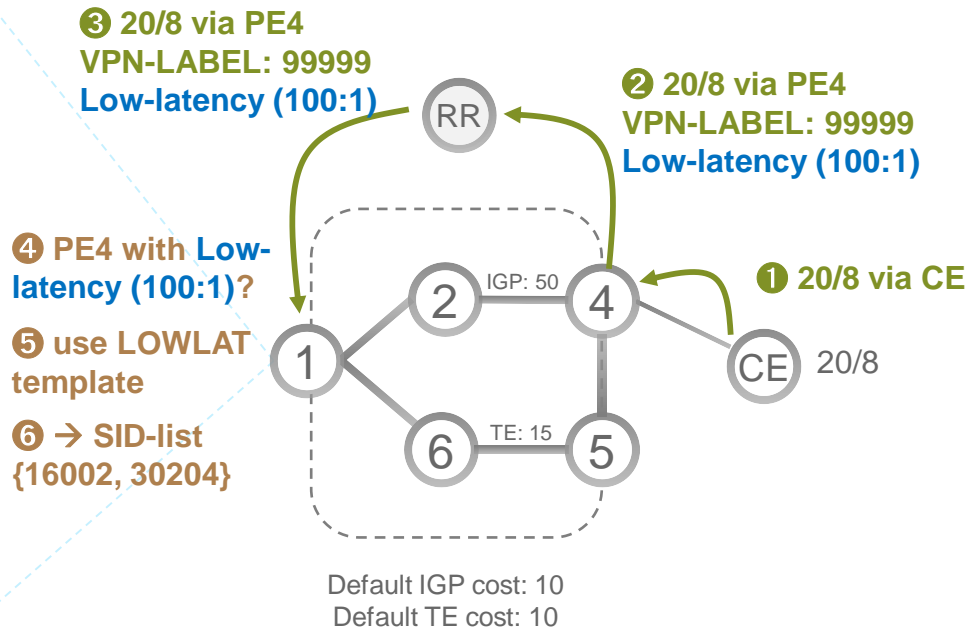
# On-demand SR Policy work-flow

5

```

route-policy ON_DEMAND_SR
  if community matches-any (100:1) then
    set mpls traffic-eng attributeset LOWLAT
  endif
  pass
end-policy
!
router bgp 1
  neighbor 1.1.1.10
  address-family vpnv4 unicast
  route-policy ON_DEMAND_SR in
!
segment-routing
  traffic-eng
  attribute-set LOWLAT
  metric
  type te

```



# Automated performant steering

8

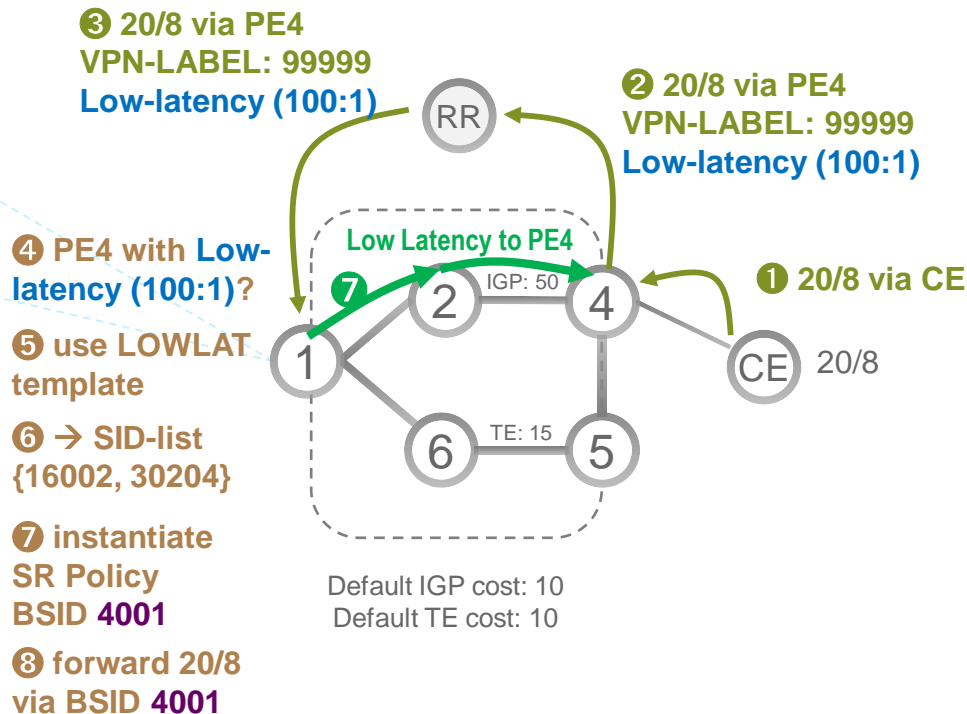
## FIB table at PE1

BGP: 20/8 via 4001  
 SRTE: 4001: Push {16002, 30204}

Automatically, the service route resolves on the Binding SID (4001) of the SR Policy it requires

Simplicity and Performance

No complex PBR to configure, no PBR performance tax

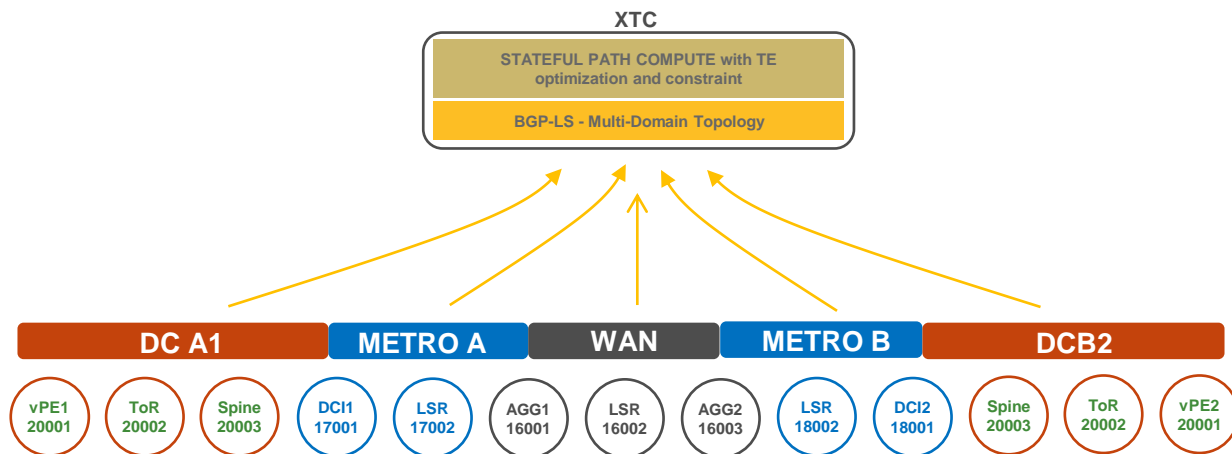


# Benefits

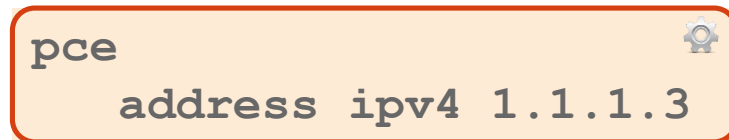
- SLA-aware BGP service
- No a-priori full-mesh of SR policy configuration
  - 3 to 4 common optimization templates are used throughout the network
    - > color => optimization objective
- No complex steering configuration
  - Automated steering of BGP routes on the right SLA path
  - Data plane performant
  - BGP PIC FRR data plane protection is preserved
  - BGP NHT fast control plane convergence is preserved

# XTC

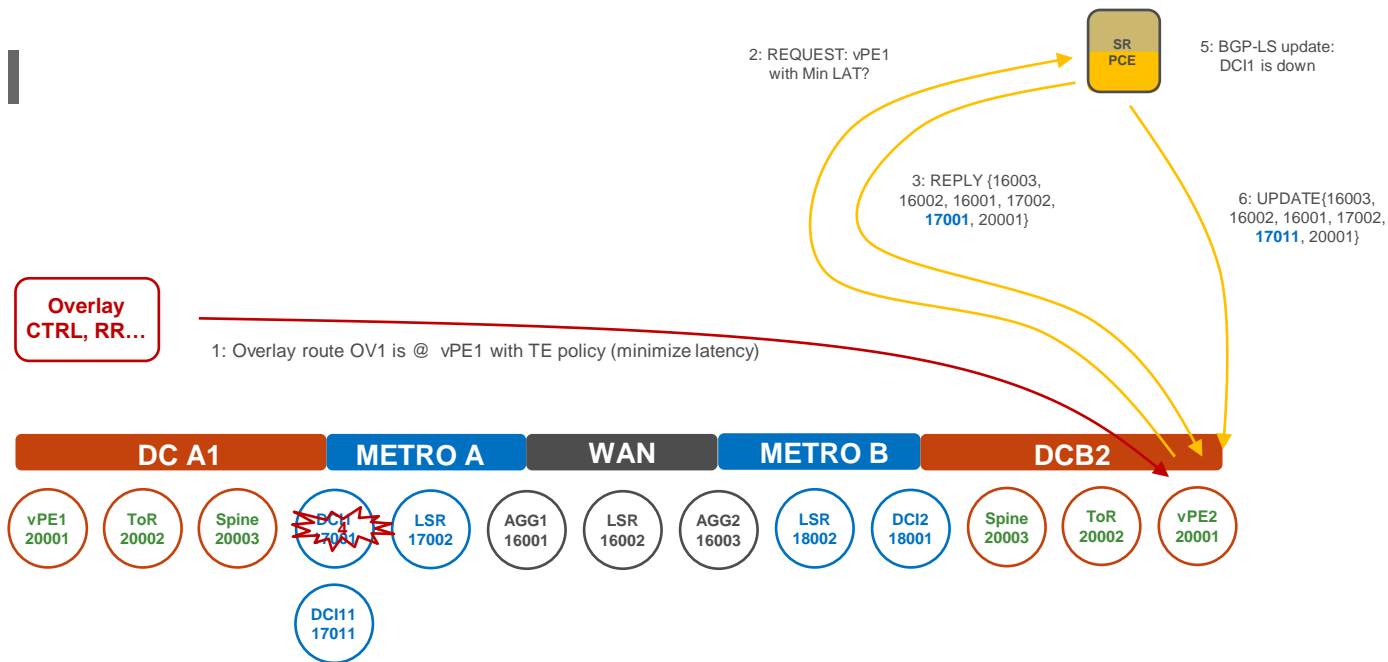
# XTC



- Multi-domain topology
  - Realtime reactive feed via BGP-LS
- Multi-domain path compute with TE optimization and constraint
  - SRTE algorithms

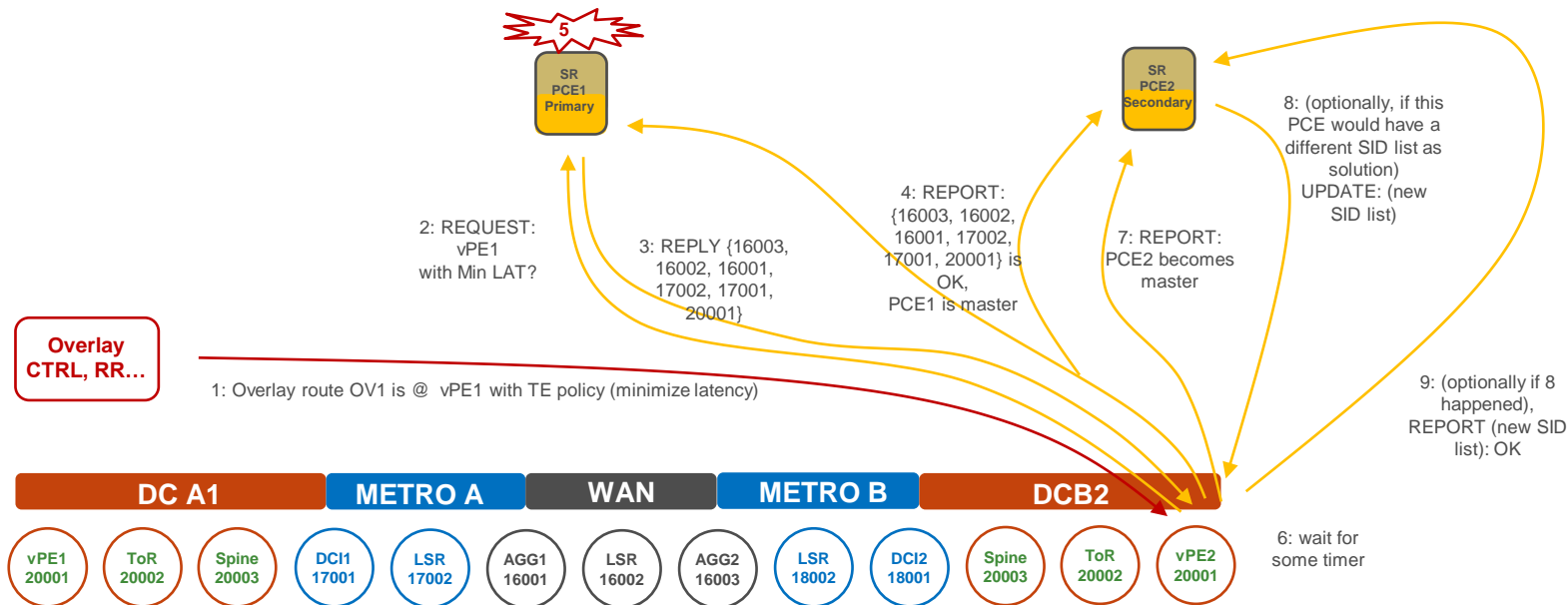


# Stateful



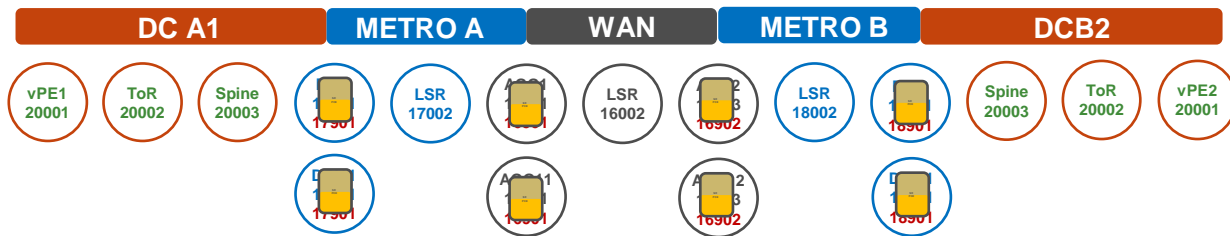
- XTC remembers the request and updates the SID list upon any topology change
  - Anycast SID's and Local FRR (TILFA) minimize traffic loss during the stateful re-optimization

# HA



- We leverage well-known standardized PCE HA

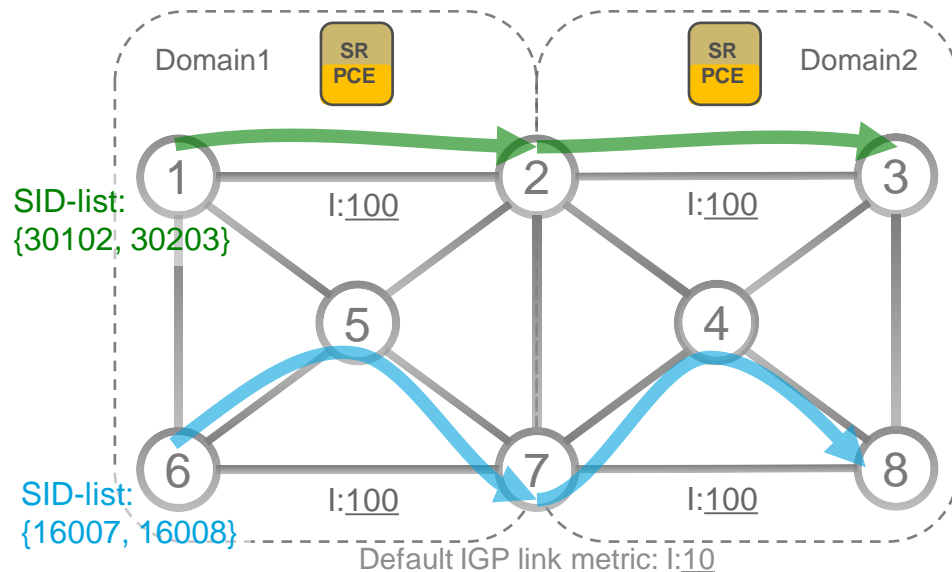
# Fundamentally Distributed



- XTC not to be considered as a single “god” box
- XTC is closer to RR
- Different vPE’s can use different pairs of XTC’s
- XTC preference can either be based on proximity or service



# Service Disjointness



```
segment-routing
traffic-eng
policy POLICY1
end-point ipv4 1.1.1.3 color 20
path
  preference 50
  dynamic mpls pce
  metric
  type igp
  association group 1 type node
```

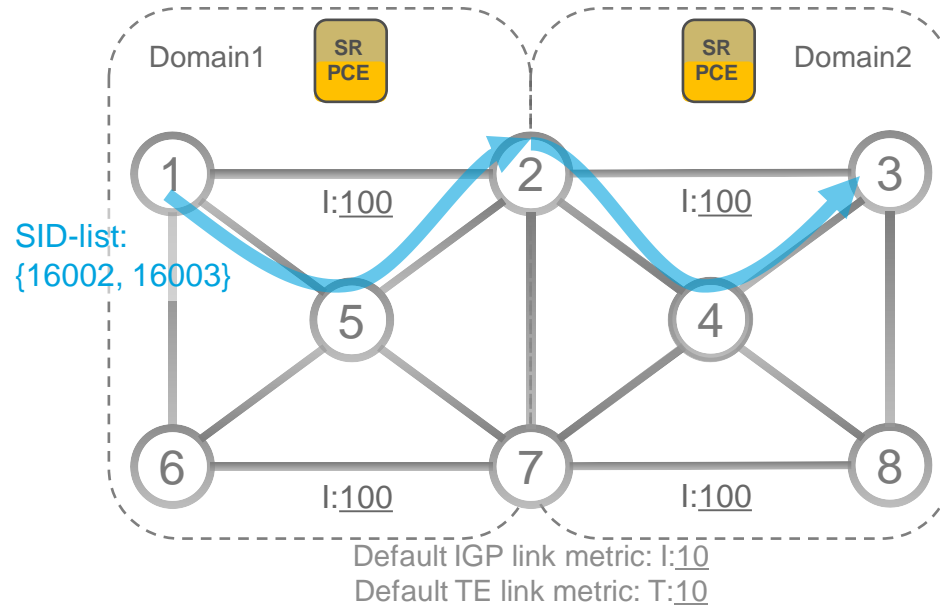
Node1

```
segment-routing
traffic-eng
policy POLICY2
end-point ipv4 1.1.1.8 color 20
path
  preference 50
  dynamic mpls pce
  metric
  type igp
  association group 1 type node
```

Node6

- Two dynamic paths between two different pairs of (headend, endpoint) must be disjoint from each other

# Inter-Domain Path – Best Effort

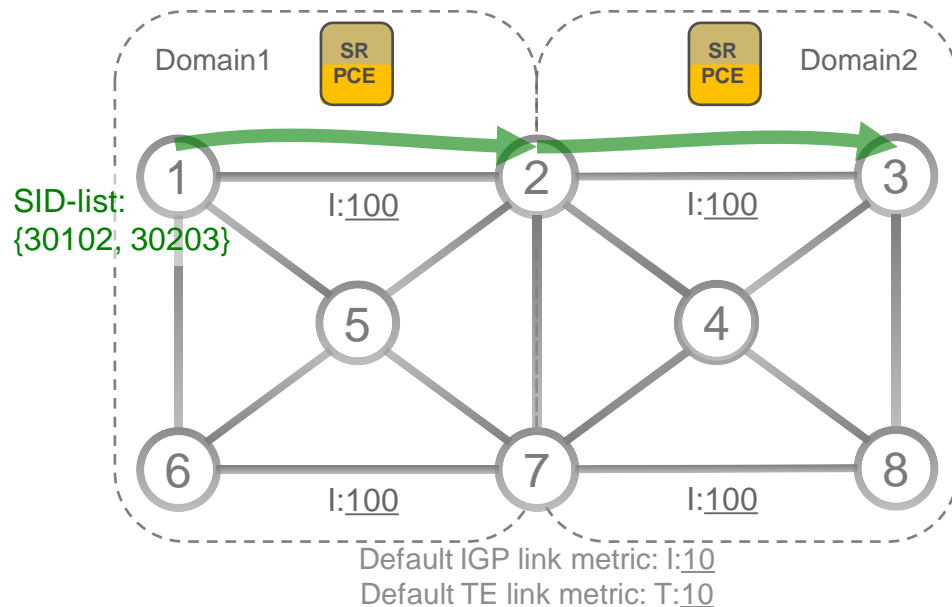


```
segment-routing
traffic-eng
policy POLICY1
end-point ipv4 1.1.1.3 color 20
path
  preference 50
  dynamic mpls pce
  metric
  type igp
```

Node1

- There is no a-priori route distribution between domains

# Inter-Domain Path – Low-Latency



```
segment-routing
traffic-eng
policy POLICY1
end-point ipv4 1.1.1.3 color 20
path
  preference 50
  dynamic mpls pce
  metric
  type te
```

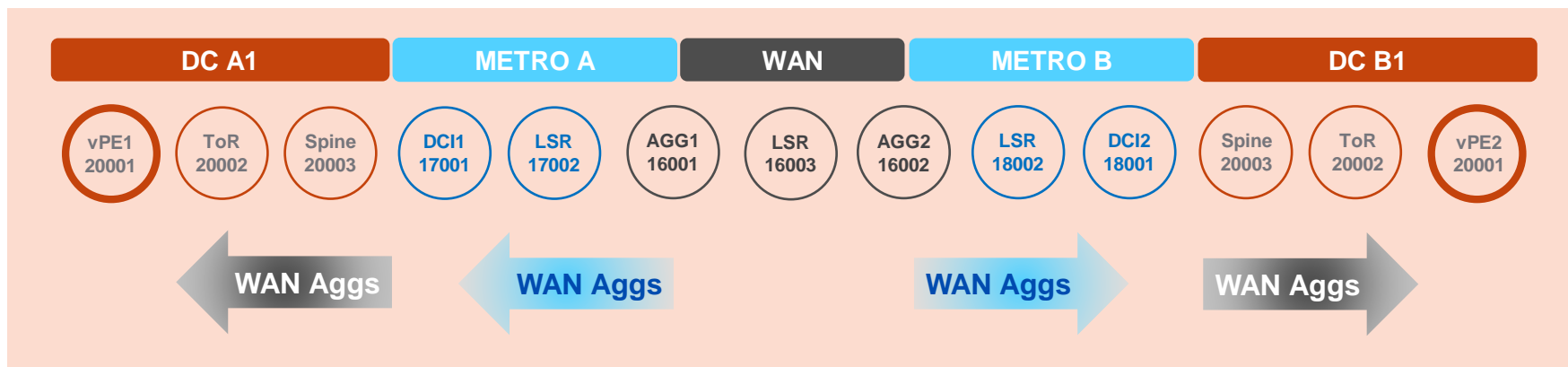
Node1

- There is no a-priori route distribution between domains
- An end-to-end policy is requested

# On-demand Next-hop

Inter-Domain

# Inter-Domain Routing

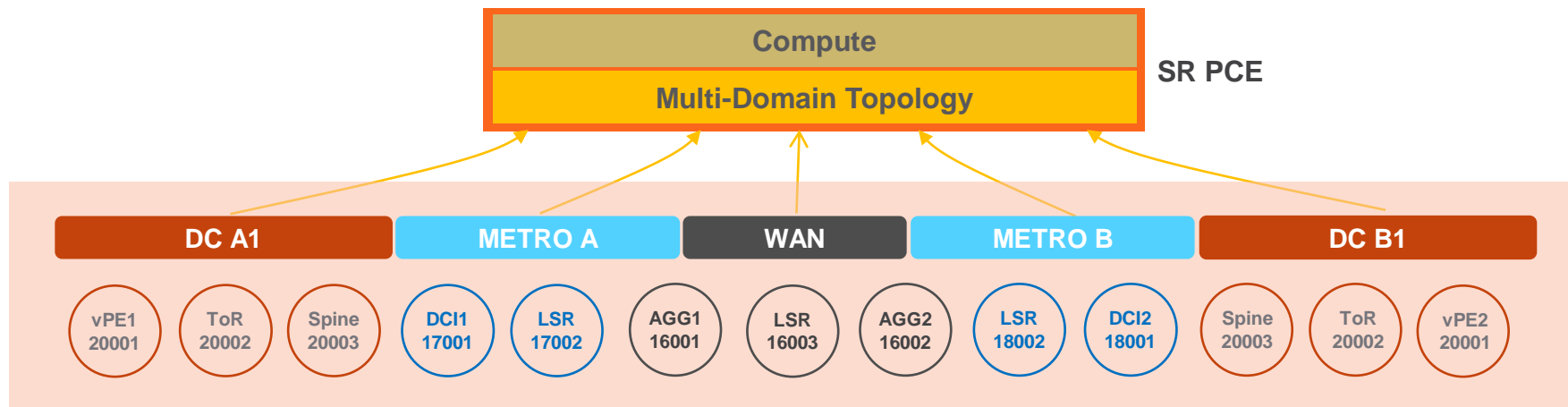


- WAN Aggs are re-distributed down to Metro and DC routing areas
- Nothing is redistributed up



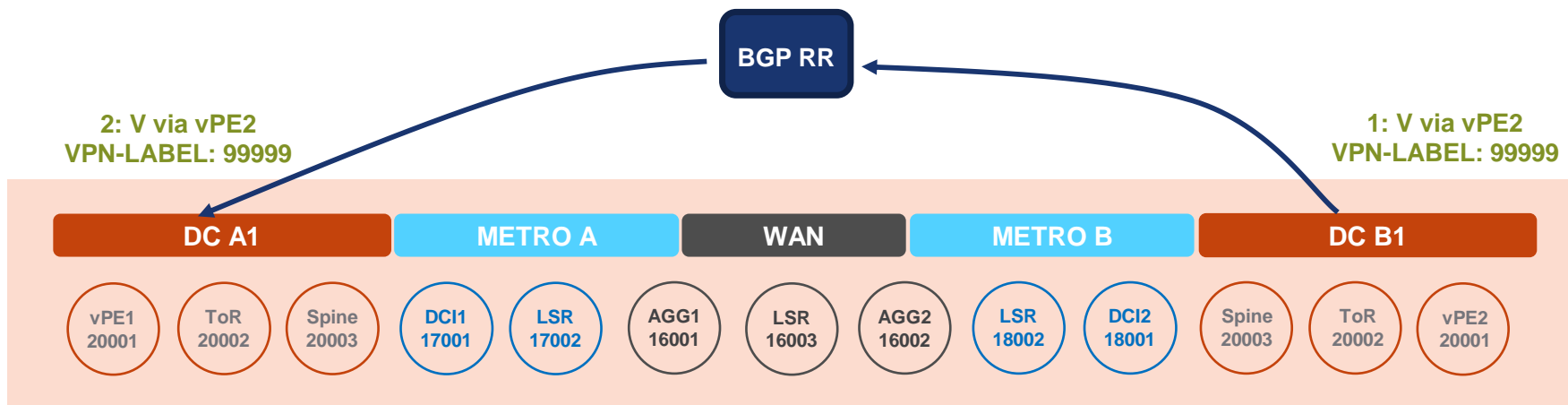
How does vPE1 reaches vPE2?

# SR Path Compute Element (PCE)



- **Multi-Domain topology**
  - ✓ **Real-time** reactive feed via BGP-LS/ISIS/OSPF from **multiple domains**
  - ✓ Including IP address and SID
- **Compute**
  - ✓ **Stateful** with **native** SRTE algorithms

# Service Provisioning



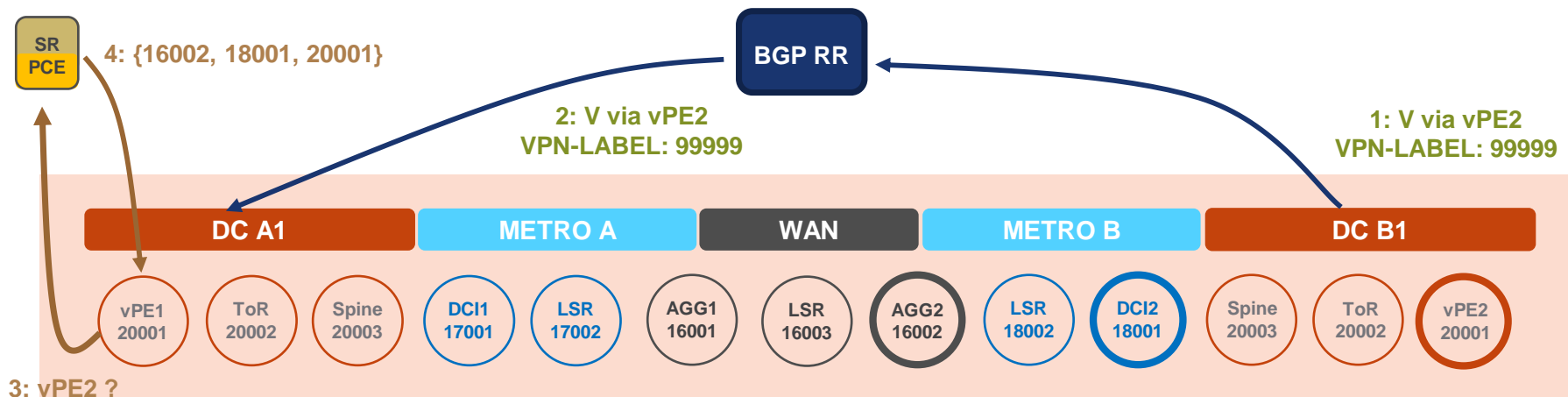
- vPE1 learns about a service route with next-hop vPE2



How does vPE1 reach the next-hop?

- ✓ vPE1 only has routes within DC A1 and to the AGG's of the WAN domain
- ✓ Solution: On-Demand Next Hop

# On-Demand SR Next-Hop Reachability

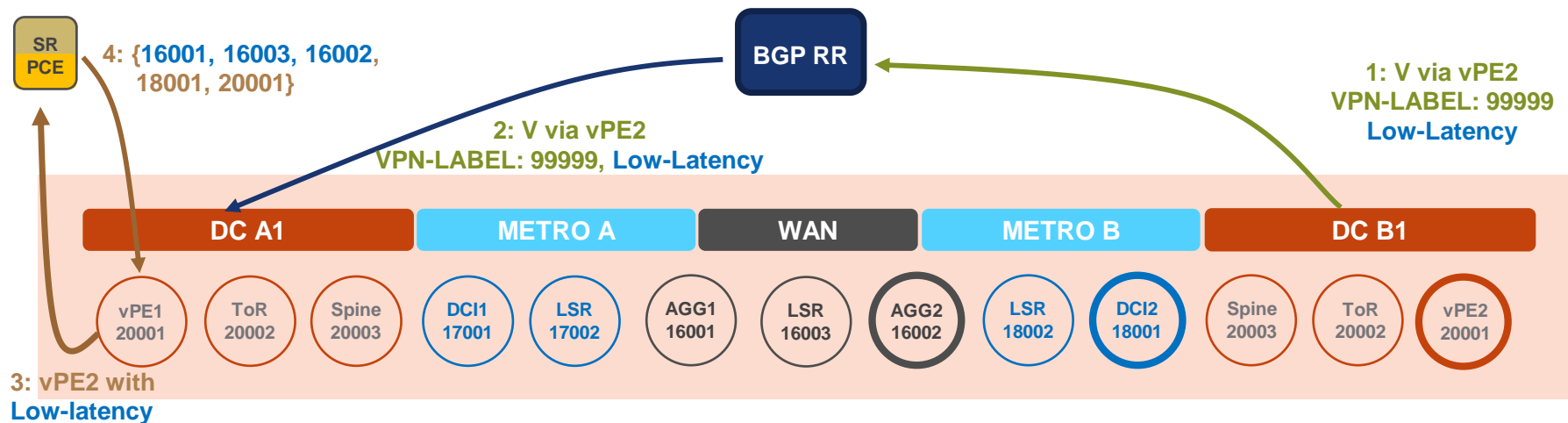


- vPE1's ODN functionality automatically requests a solution from SR-PCE
- **Scalable** - vPE1 only gets the inter-domain paths that it needs
- **Simple** - no BGP3107 pushing all routes everywhere



# On-Demand SR Next-Hop

## End-to-End Policy



- vPE1's ODN functionality automatically requests a solution from SR-PCE
- **Scalable** - vPE1 only gets the inter-domain paths that it needs
- **Simple** - no BGP3107 pushing all routes everywhere

# ODN config at PE1

```
route-policy ON_DEMAND_SR
  if community matches-any (100:1) then
    set mpls traffic-eng attributeset LOWLAT
  endif
  pass
end-policy
!
router bgp 1
  neighbor 1.1.1.10
  address-family vpnv4 unicast
  route-policy ON_DEMAND_SR in
!
segment-routing
  traffic-eng
  attribute-set LOWLAT
  pce
  metric
  type te
```

# Conclusion

# SR TE

- Simple, Automated and Scalable
  - No core state: **state in the packet header**
  - No tunnel interface: “**SR Policy**”
  - No headend a-priori configuration: **on-demand** policy instantiation
  - No headend a-priori steering: **on-demand** steering
- Multi-Domain
  - **XTC**
- Lots of Functionality
  - Designed with **lead operators** along their use-cases

# Your Time Is Now

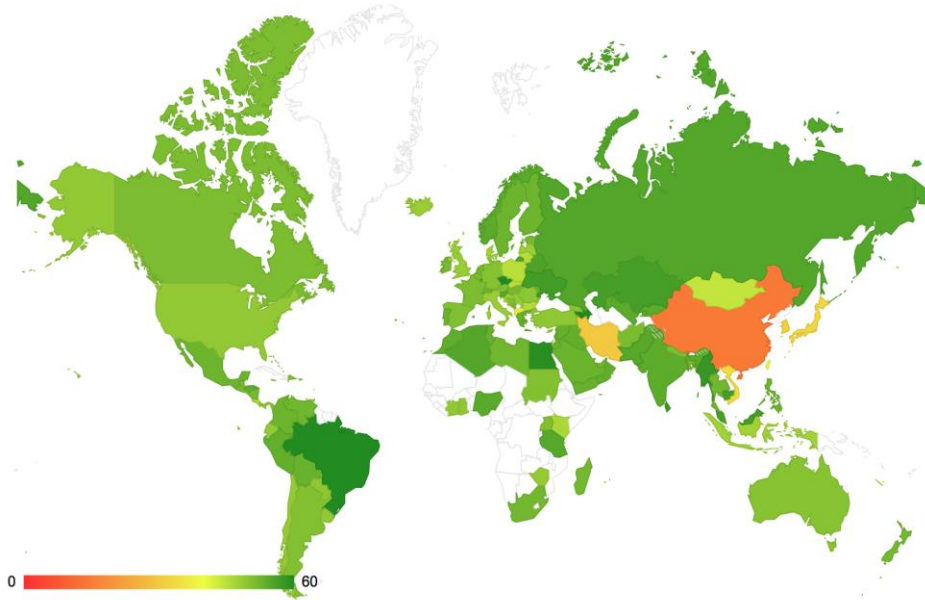
# SRv6

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BRKRST-3122



# IPv6 adoption is a reality



% Web pages available over IPv6

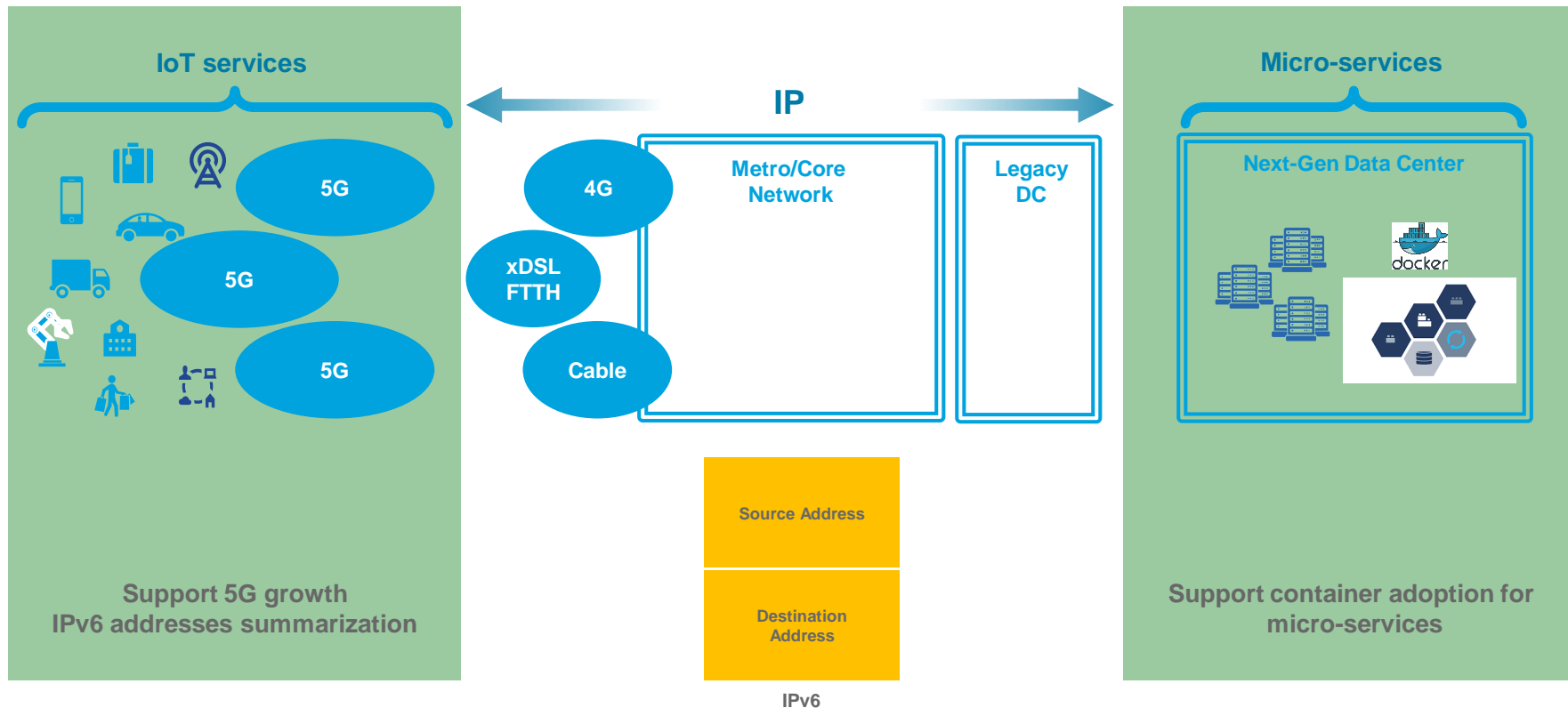
Source: [6lab.cisco.com](http://6lab.cisco.com) – World maps – Web content

Global IPv6 traffic  
grew 243% in 2015

Globally IPv6 traffic will  
grow 16-fold from 2015 to  
2020

IPv6 will be 34% of total  
Internet traffic in 2020

# IPv6 provides reachability





# SRv6 – Segment Routing & IPv6

SRv6 for anything else

IPv6 for reach

- Simplicity
  - Protocol elimination
- SLA
  - FRR and TE
- Overlay
- NFV
- SDN
  - SR is de-facto SDN architecture
- 5G Slicing

# SRv6 for underlay

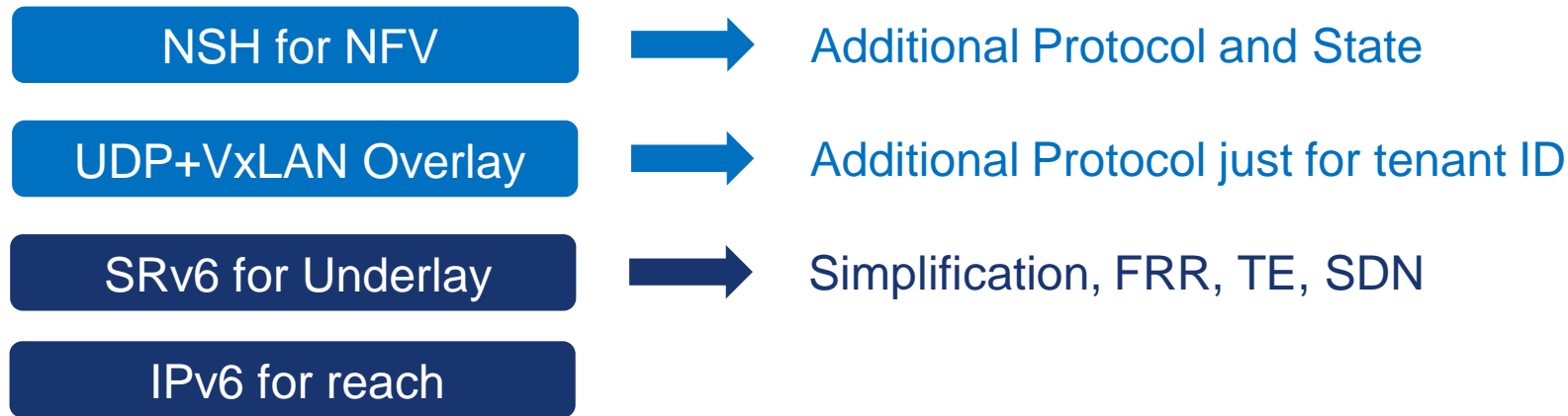
RSVP for FRR/TE

IPv6 for reach



Horrendous states scaling in  $k \cdot N^2$

# Opportunity for further simplification



- Multiplicity of protocols and states hinder network economics

# Our commitment to Lead Operators



**Standardization**  
IETF

**Seamless  
Deployment**

**Multi-vendor  
Consensus**

- Clear track record for SR team



# SR for anything Network as a Computer

# Network instruction

Locator

Function

- 128-bit SRv6 SID
  - Locator: routed to the node performing the function
  - Function: any possible function (optional argument)  
either local to NPU or app in VM/Container
  - Flexible bit-length selection

# Network instruction

Locator

Function(arg)

- 128-bit SRv6 SID
  - Locator: routed to the node performing the function
  - Function: any possible function (optional argument)  
either local to NPU or app in VM/Container
  - Flexible bit-length selection

# Network instruction

Locator

Function

- 128-bit SRv6 SID
  - Locator: routed to the node performing the function
  - Function: any possible function (optional argument)  
either local to NPU or app in VM/Container
  - Flexible bit-length selection



# Network Program

Next Segment



Locator 1

Function 1

Locator 2

Function 2

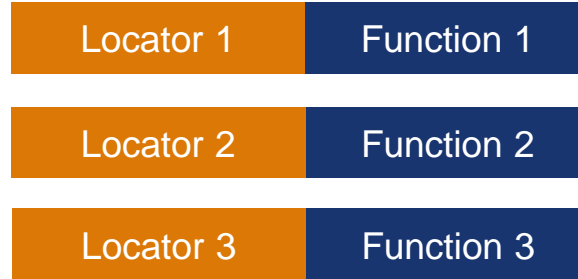
Locator 3

Function 3



# Network Program

Next Segment →



# Network Program

Next Segment



Locator 1

Function 1

Locator 2

Function 2

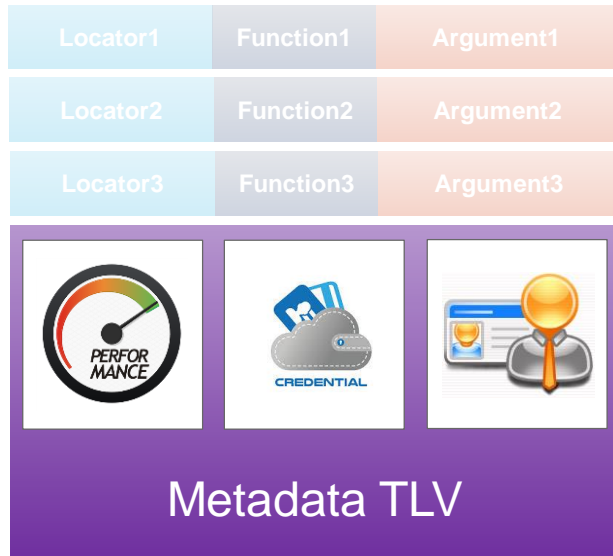
Locator 3

Function 3

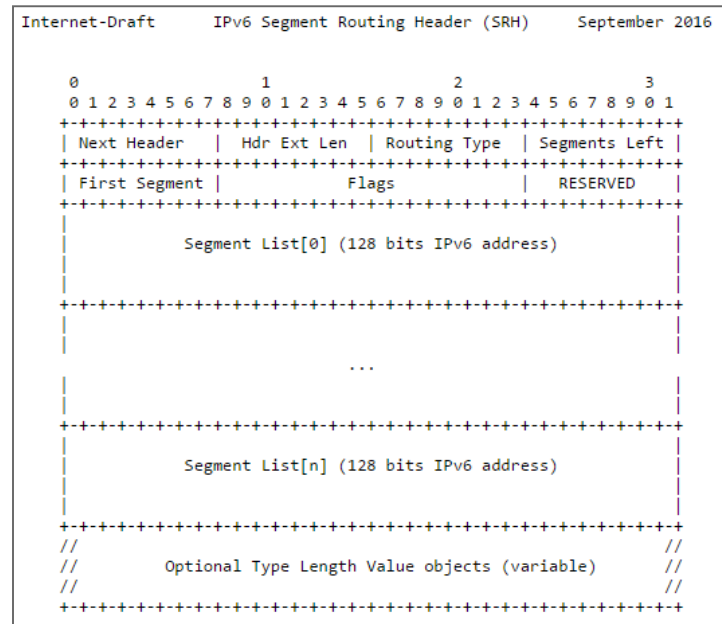
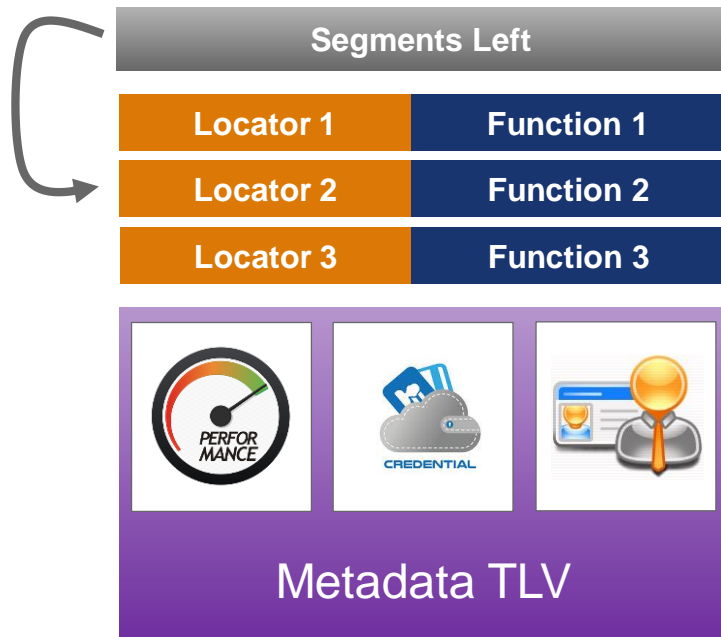


# Argument shared between functions

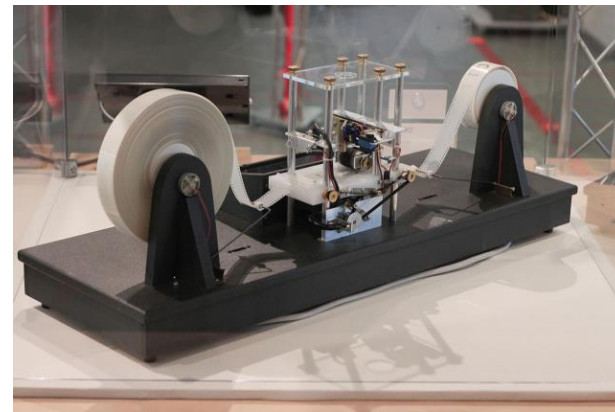
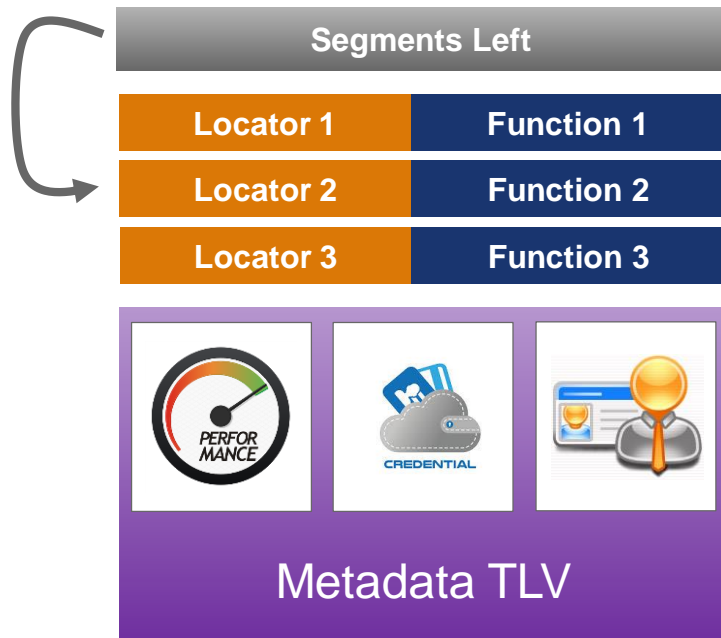
**“Global”  
Argument**



**Cisco** *live!*

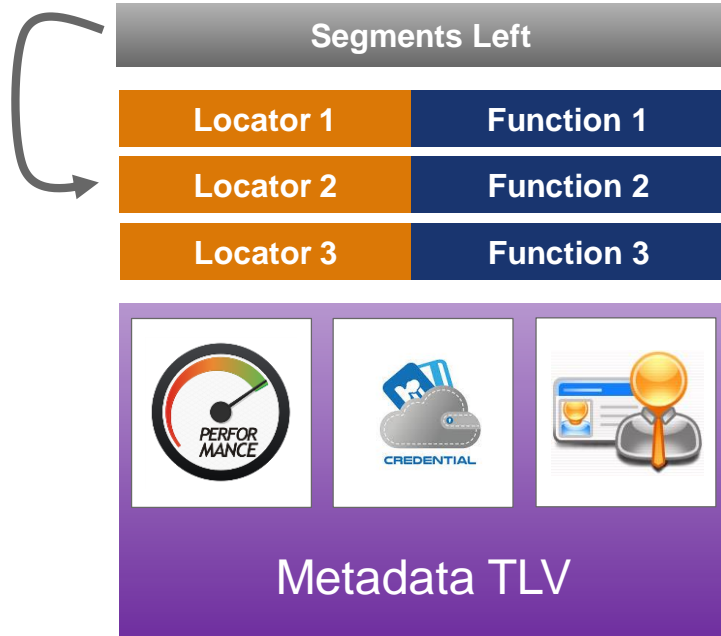


# SRv6 for anything



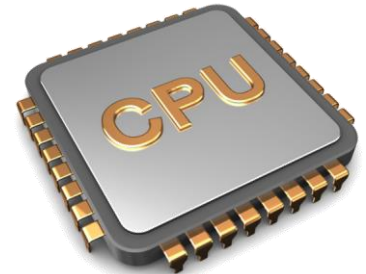
## Turing

# SRv6 for anything



Optimized for HW processing  
e.g. Underlay & Tenant use-cases

Optimized for SW processing  
e.g. NFV, Container, Micro-Service



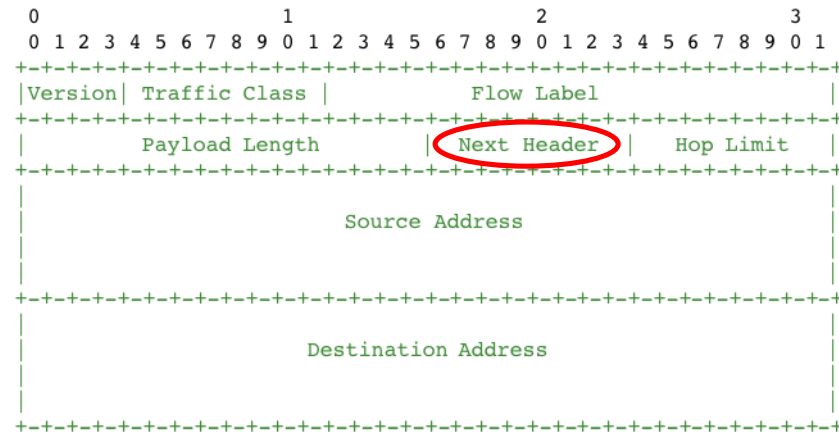


# SR Header

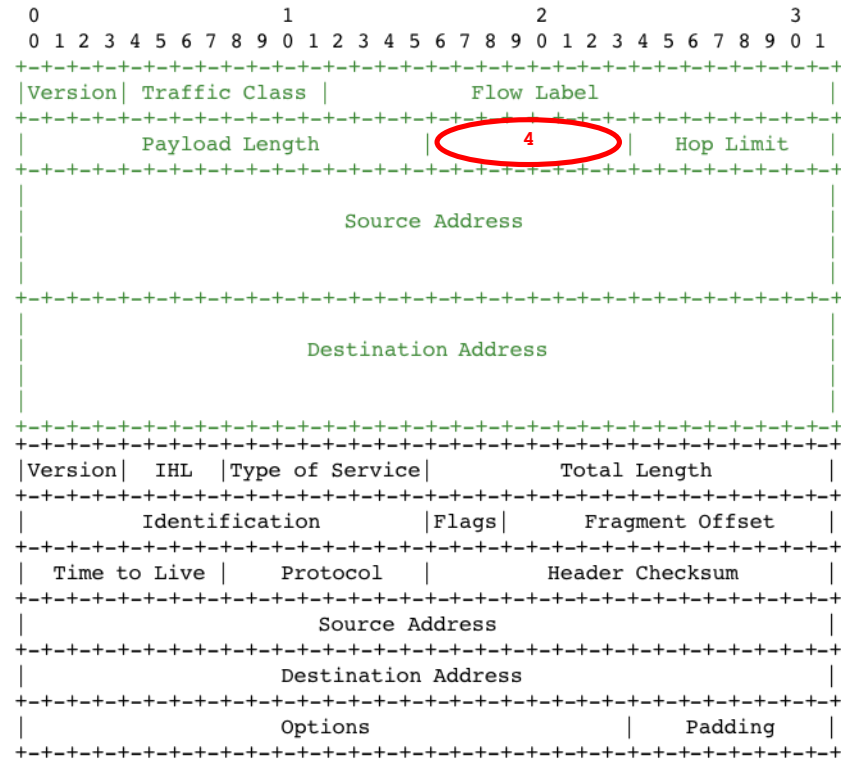


# IPv6 Header

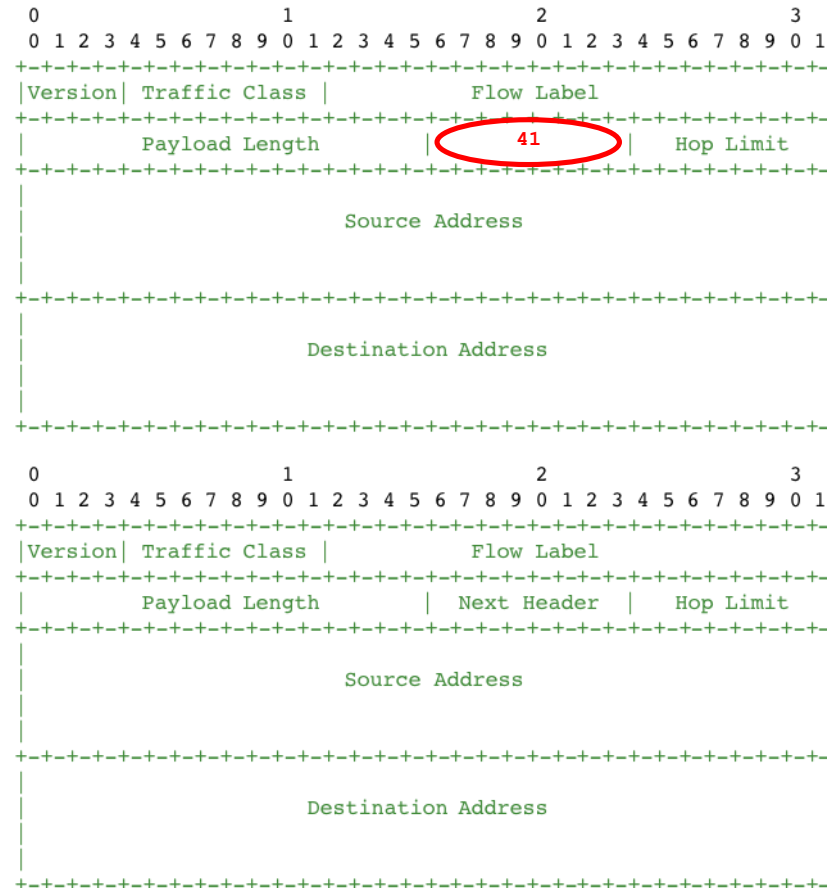
- Next Header (NH)
  - Indicates what comes next



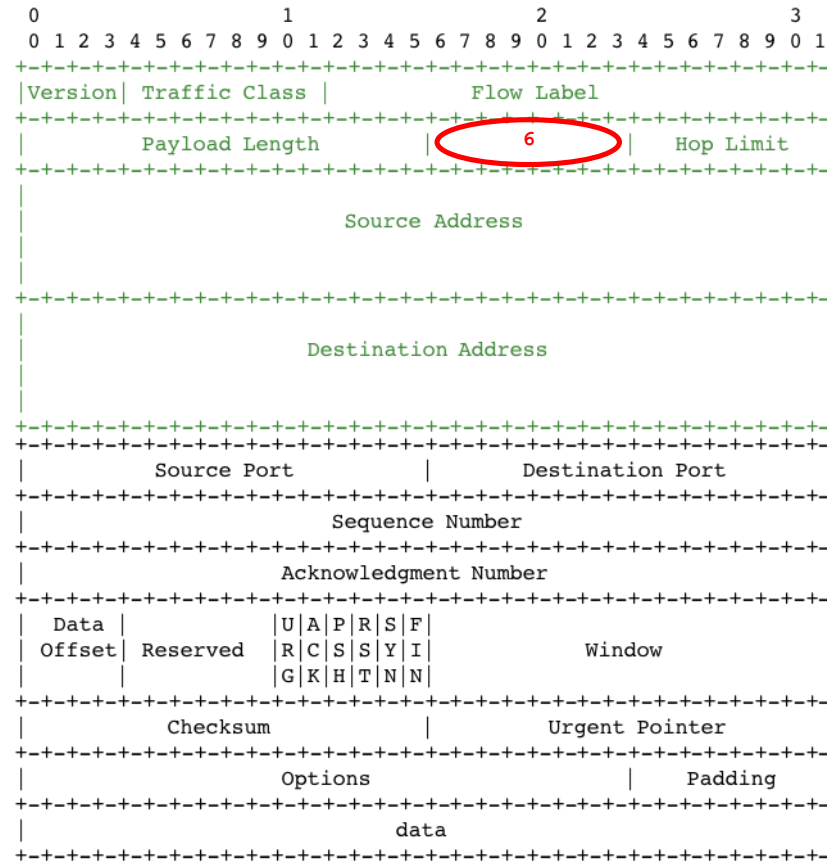
# NH = IPv4



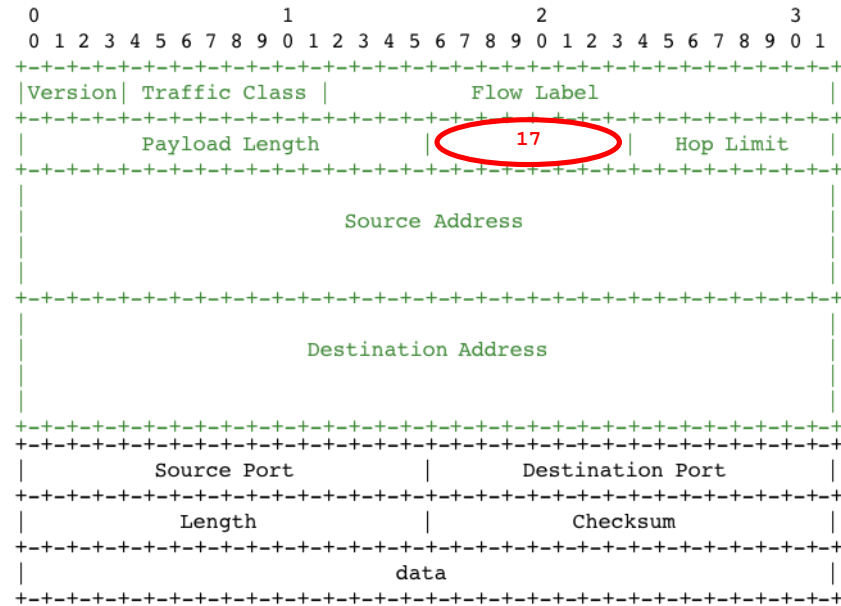
# NH = IPv6



NH = TCP

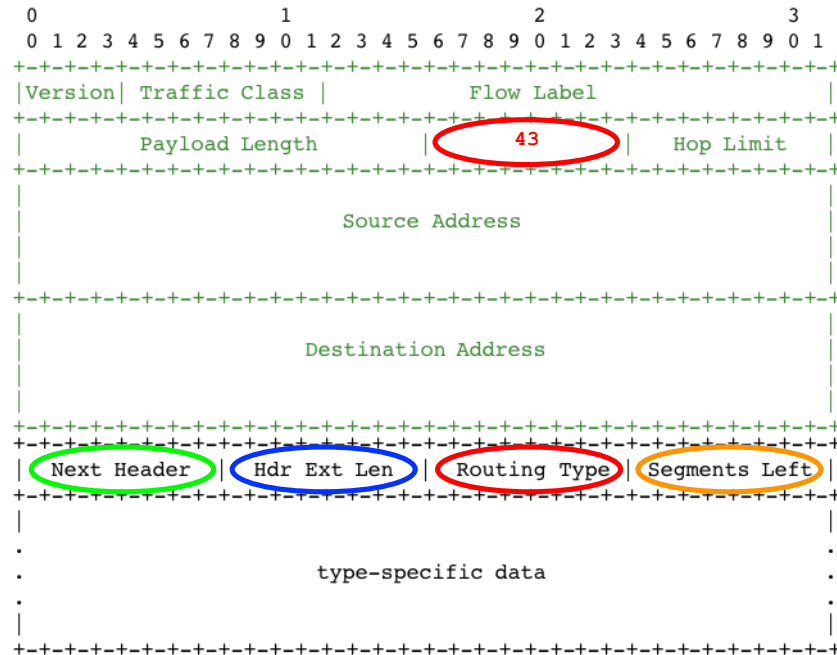


# NH = UDP



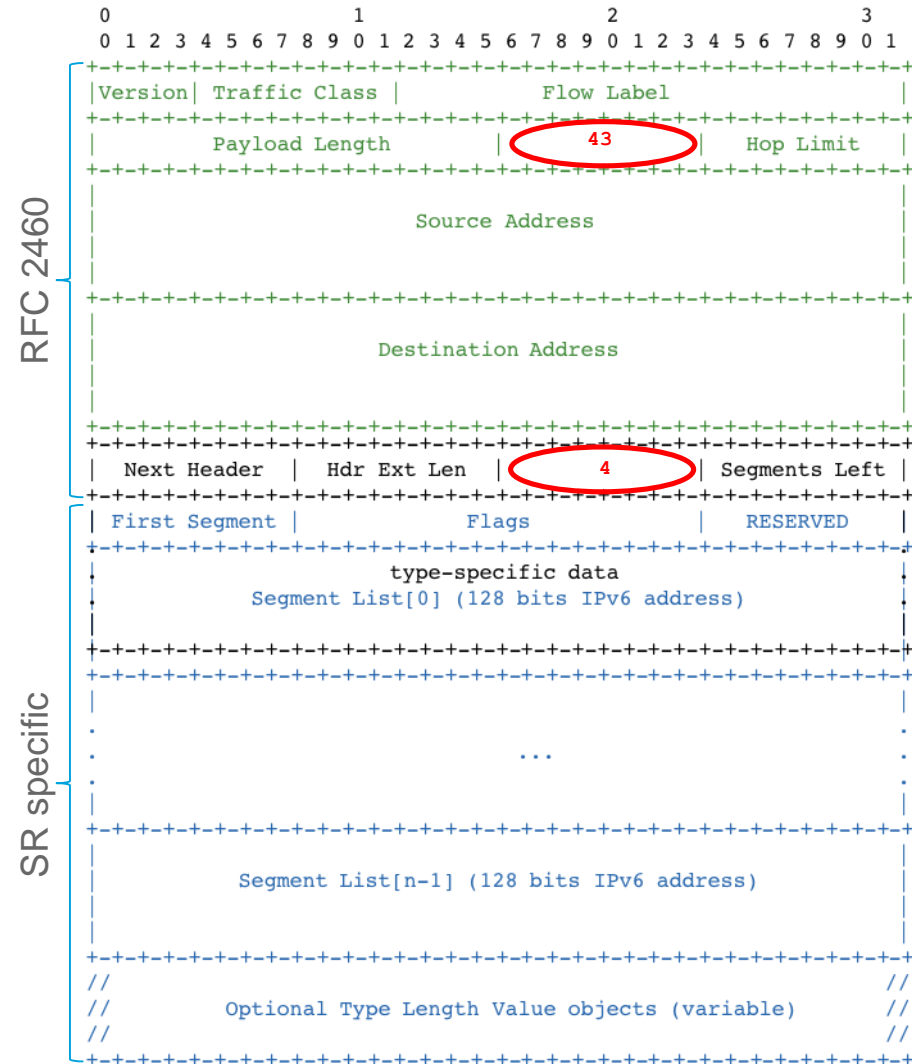
# NH = Routing Extension

- Generic routing extension header
  - Defined in RFC 2460
  - Next Header: UDP, TCP, IPv6...
  - Hdr Ext Len: **Any IPv6 device can skip this header**
  - Segments Left: **Ignore extension header if equal to 0**
- Routing Type field:
  - > 0 Source Route (deprecated since 2007)
  - > 1 Nimrod (deprecated since 2009)
  - > 2 Mobility (RFC 6275)
  - > 3 RPL Source Route (RFC 6554)
  - > 4 Segment Routing



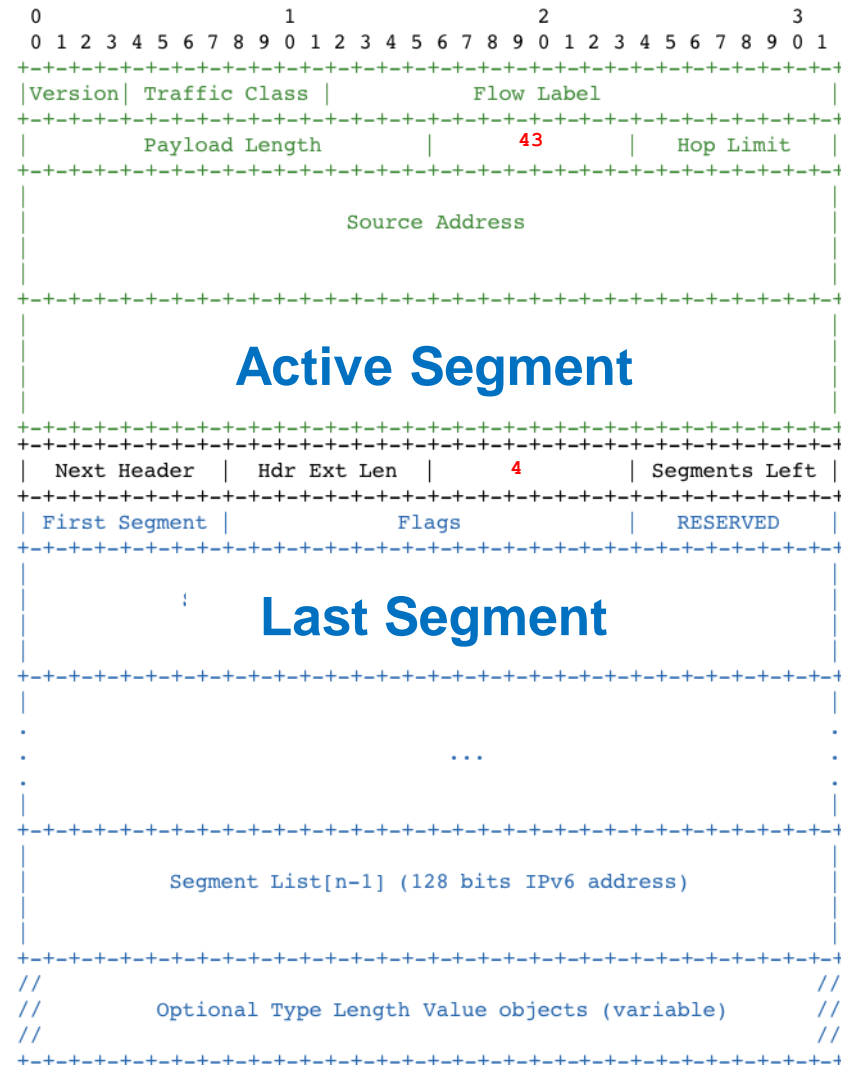
# NH = SRv6

- NH = 43, Type = 4



# SRH

- SRH contains
  - the list of segments
  - Segments left (SL)
  - Flags
  - TLV
- Active segment is in the IPv6 DA
- Next segment is at index SL-1
- The last segment is at index 0
  - Reversed order

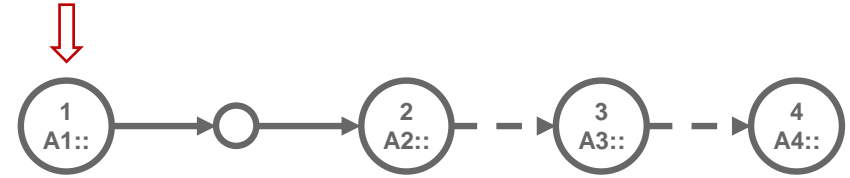




# SRH Processing

# Source Node

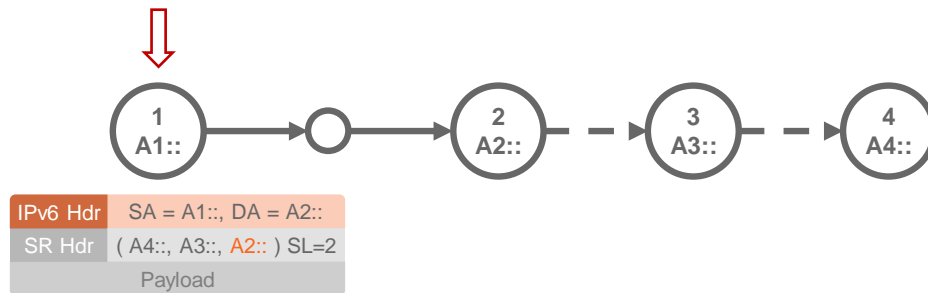
- Source node is SR-capable
- SR Header (SRH) is created with
  - Segment list in reversed order of the path
    - > Segment List [ 0 ] is the LAST segment
    - > Segment List [  $n - 1$  ] is the FIRST segment
  - Segments Left is set to  $n - 1$
  - First Segment is set to  $n - 1$
- IP DA is set to the first segment
- Packet is send according to the IP DA
  - Normal IPv6 forwarding



IPv6 Hdr	SA = A1::, DA = A2::
SR Hdr	( A4::, A3::, A2:: ) SL=2
	Payload

IPv6 Hdr	Version	Traffic Class	Flow Label	
	Payload Length		Next = 43	Hop Limit
	Source Address = A1::			
	Destination Address = A2::			
SR Hdr	Next Header	Len= 6	Type = 4	SL = 2
	First = 2	Flags		RESERVED
	Segment List [ 0 ] = A4::			
	Segment List [ 1 ] = A3::			
	Segment List [ 2 ] = A2::			
Payload				

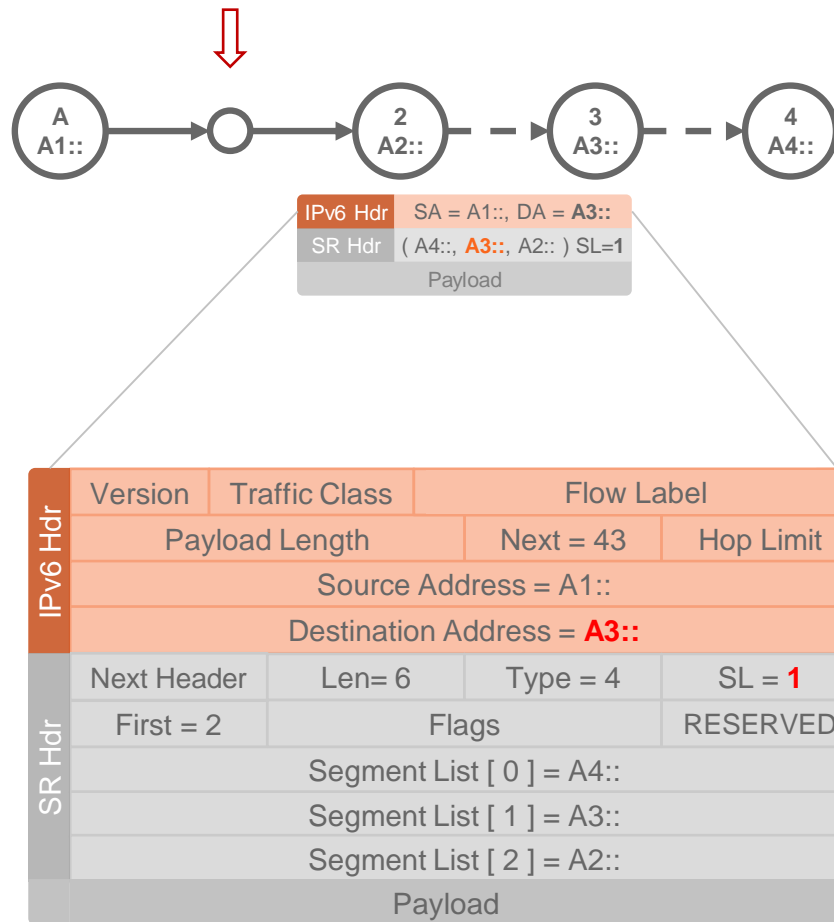
# Non-SR Transit Node



- Plain IPv6 forwarding
- Solely based on IPv6 DA
- No SRH inspection or update

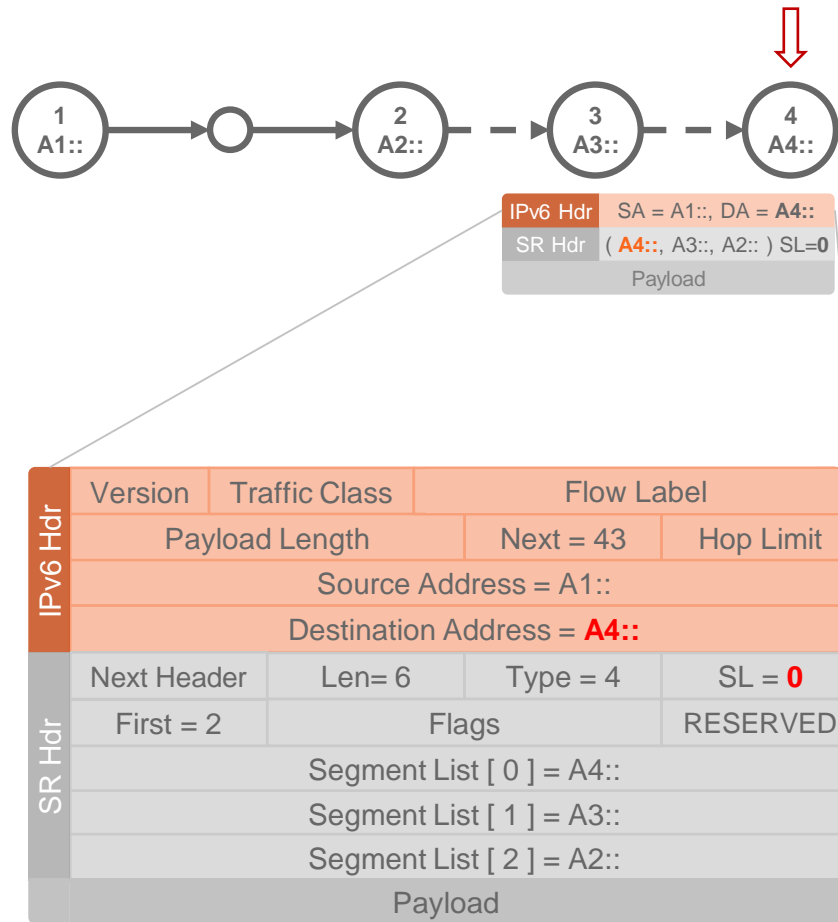
# SR Segment Endpoints

- SR Endpoints: SR-capable nodes whose address is in the IP DA
- SR Endpoints inspect the SRH and do:
  - IF Segments Left > 0, THEN
    - > Decrement Segments Left ( -1 )
    - > Update DA with Segment List [ Segments Left ]
    - > Forward according to the new IP DA



# SR Segment Endpoints

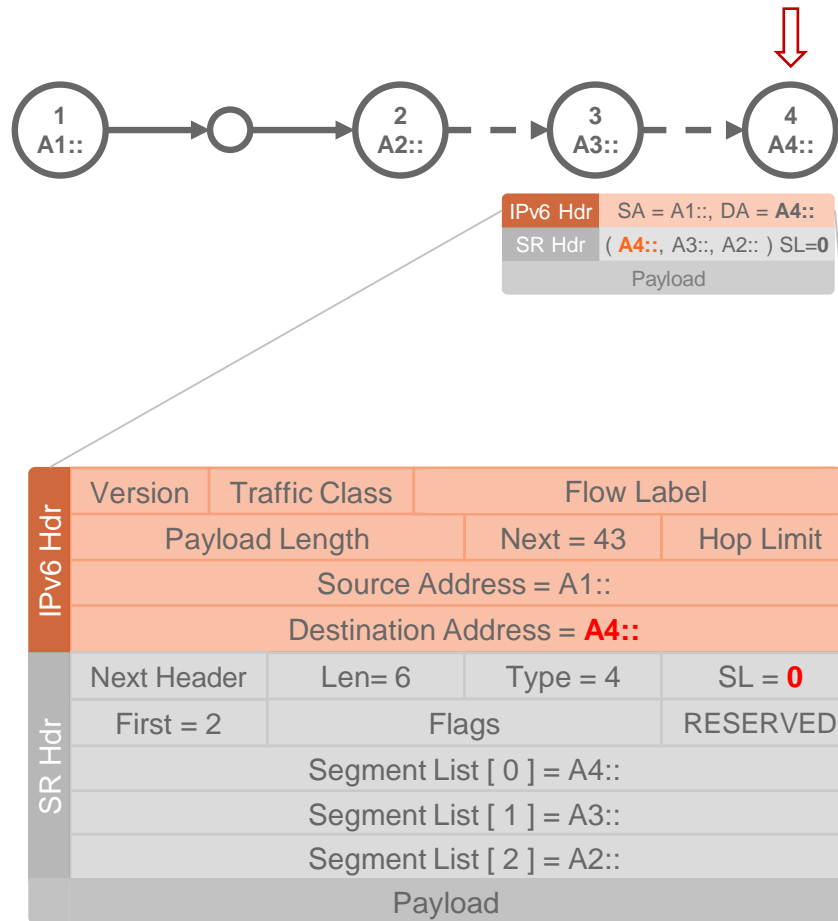
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  - ELSE (Segments Left = 0)
    - > Remove the IP and SR header
    - > Process the payload:
      - Inner IP: Lookup DA and forward
      - TCP / UDP: Send to socket
      - ...



# SR Segment Endpoints

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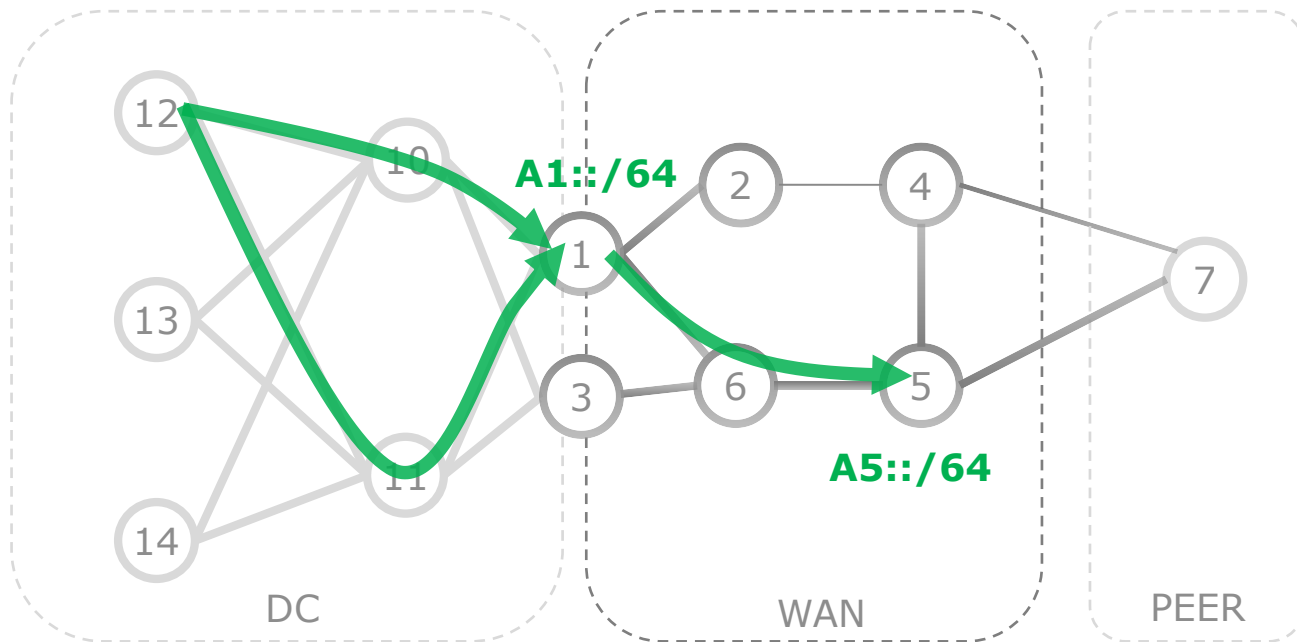
**Standard IPv6 processing**  
*The final destination does not have to be SR-capable.*



# Use-Cases

# SID allocation for illustration purpose

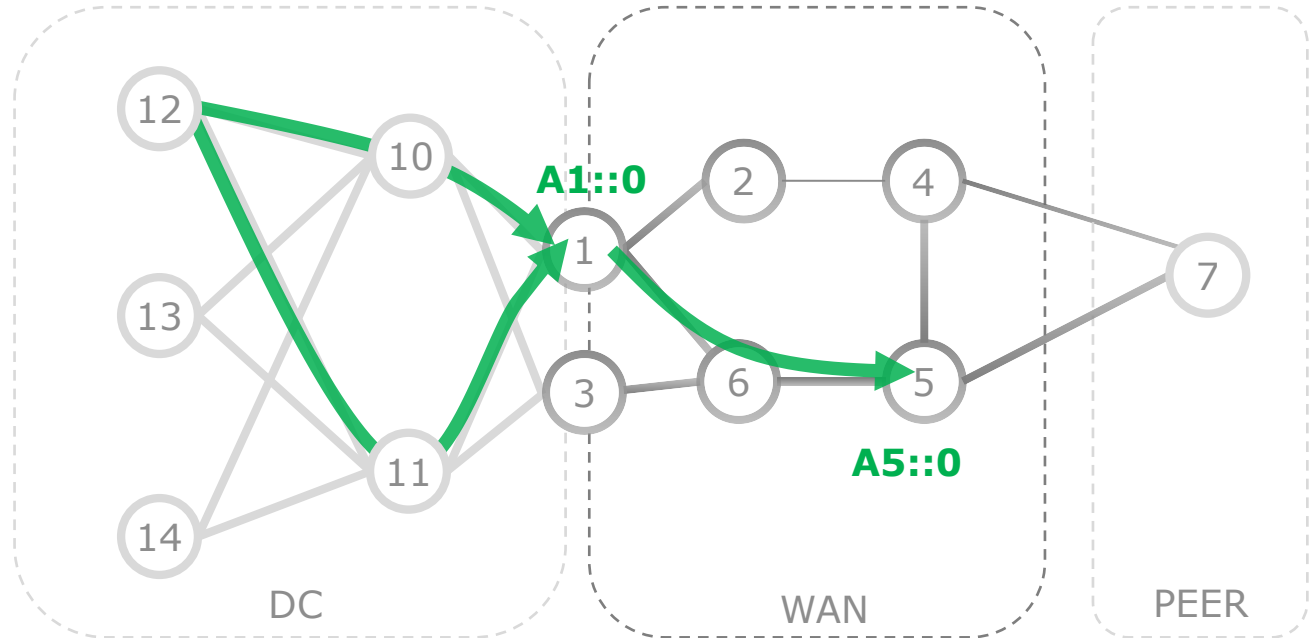
- For simplicity
- Node K advertises prefix  $AK::/64$
- The function is encoded in the last 64 bits





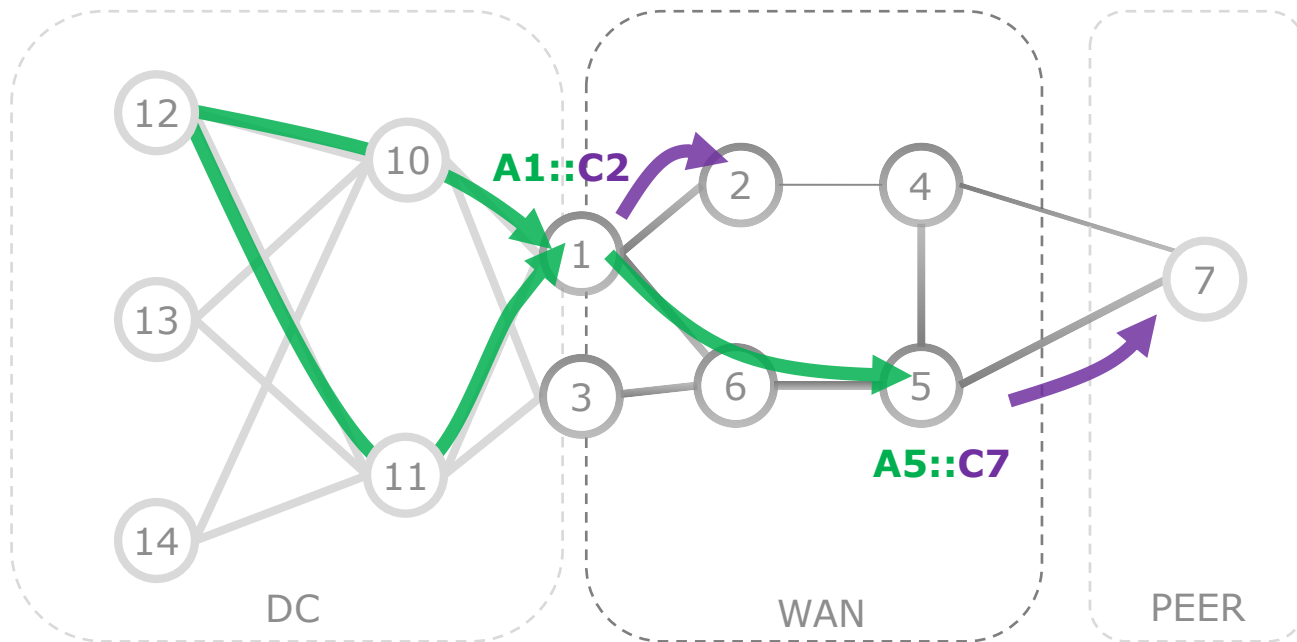
# Endpoint

- For simplicity
- Function 0 denotes the most basic function
- Shortest-path to the Node



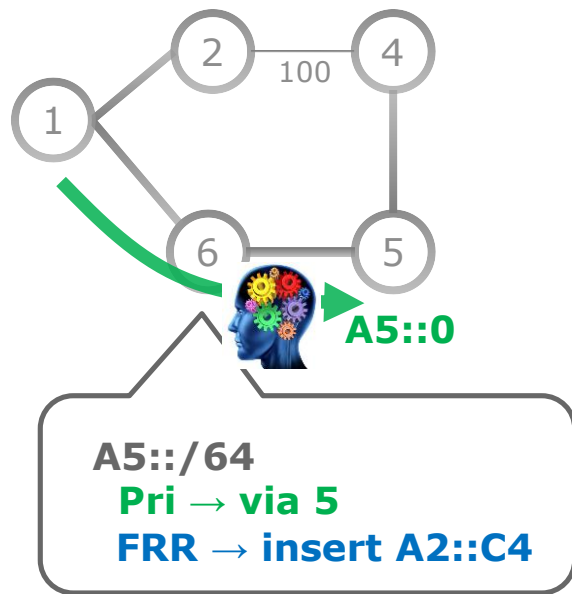
# Endpoint then xconnect to neighbor

- For simplicity
- $AK::CJ$  denotes  
Shortest-path to the  
Node K and then  
x-connect (function C)  
to the neighbor J



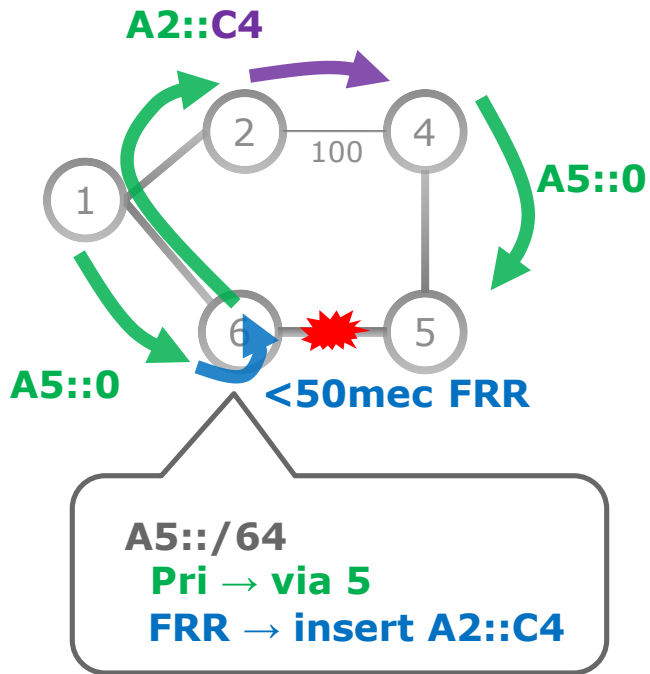
# TILFA

- **50msec Protection** upon local **link**, **node** or **SRLG** failure
- **Simple** to operate and understand
  - automatically computed by the router's IGP process
  - 100% coverage across any topology
  - predictable (backup = postconvergence)
- **Optimum backup path**
  - leverages the post-convergence path, planned to carry the traffic
  - avoid any intermediate flap via alternate path
- **Incremental deployment**
- **Distributed and Automated Intelligence**

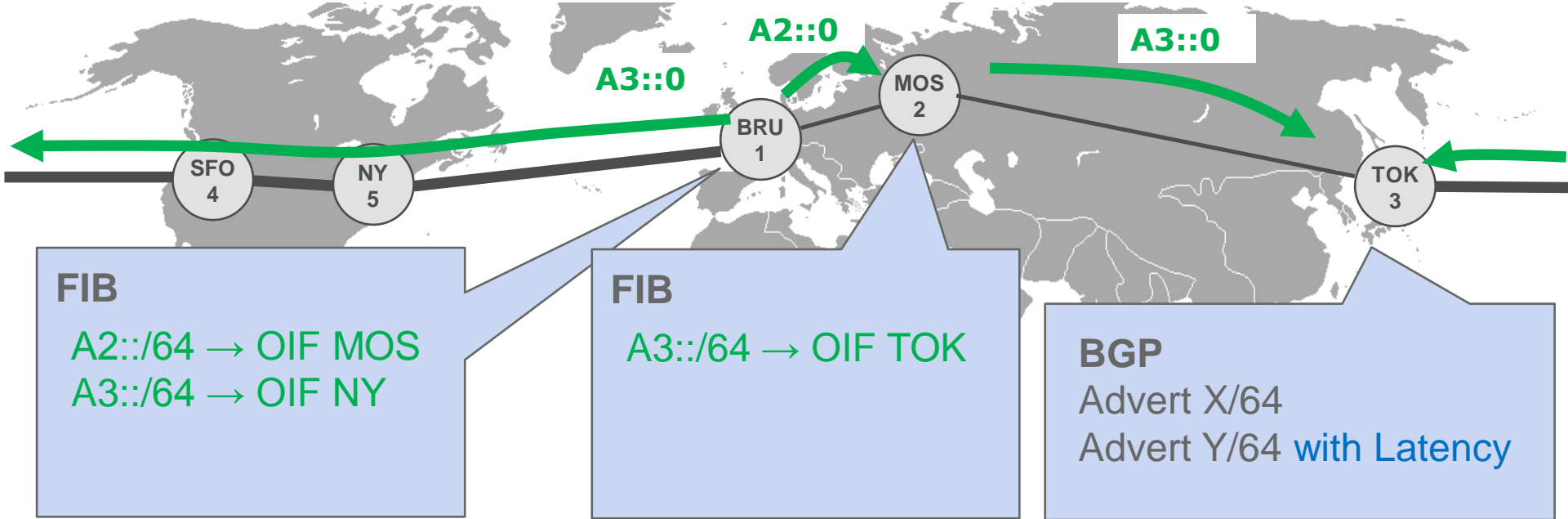


# TILFA

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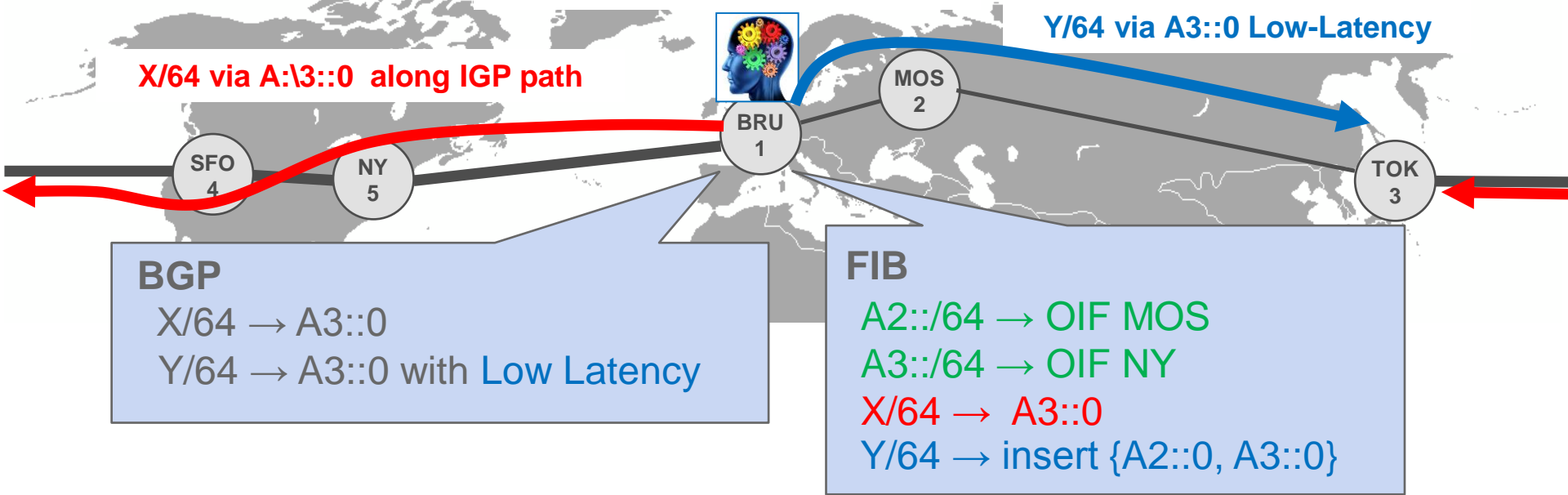
# Distributed & Automated TE



- IGP minimizes cost instead of latency

# Distributed & Automated TE

## On-Demand Distributed TE



- Distributed and Automated Intelligence
- Dynamic SRTE Policy triggered by learning a BGP route with SLA contract
- No PBR steering complexity, No PBR performance tax, No RSVP, No tunnel to configure

# Centralized TE

## Input Acquisition

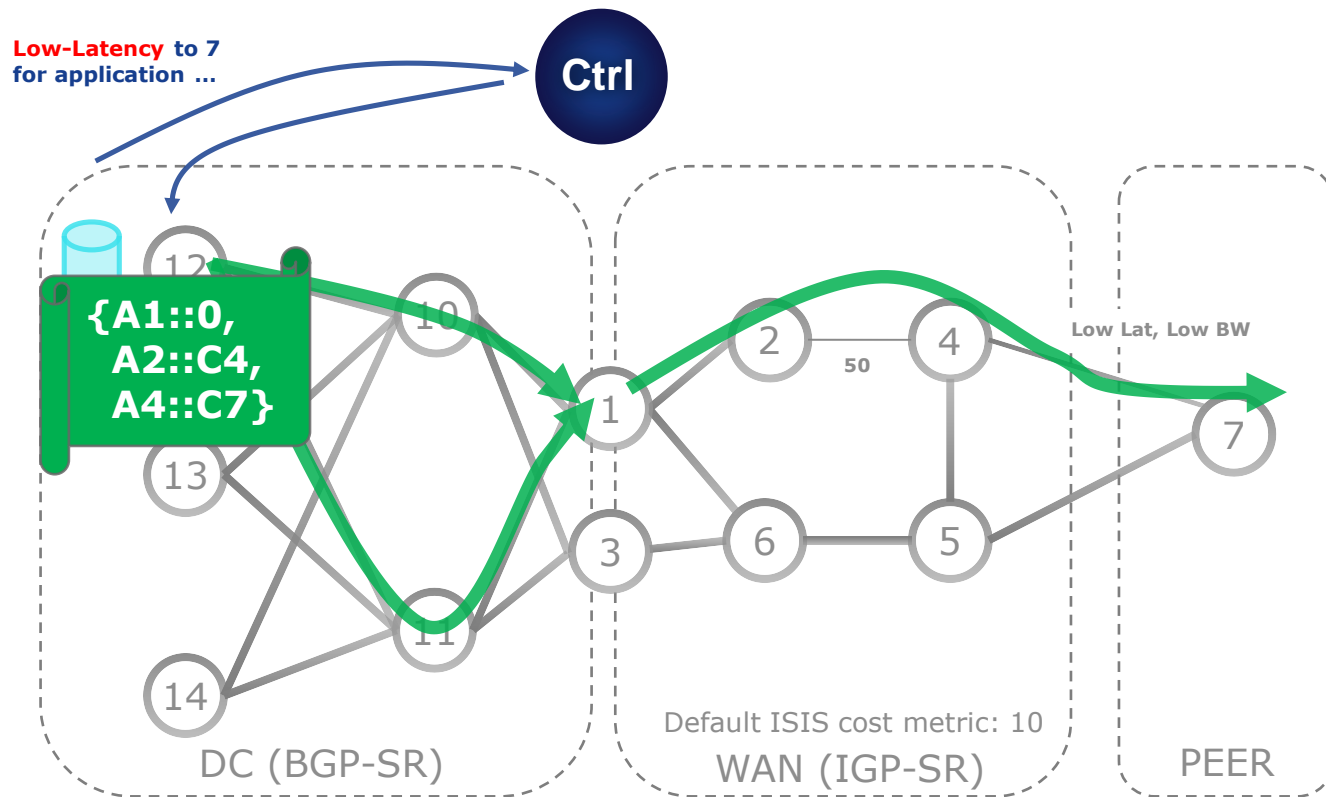
- BGP-LS
- Telemetry

## Policy Instanciation

- PCEP
- BGP-TE
- Netconfig / Yang

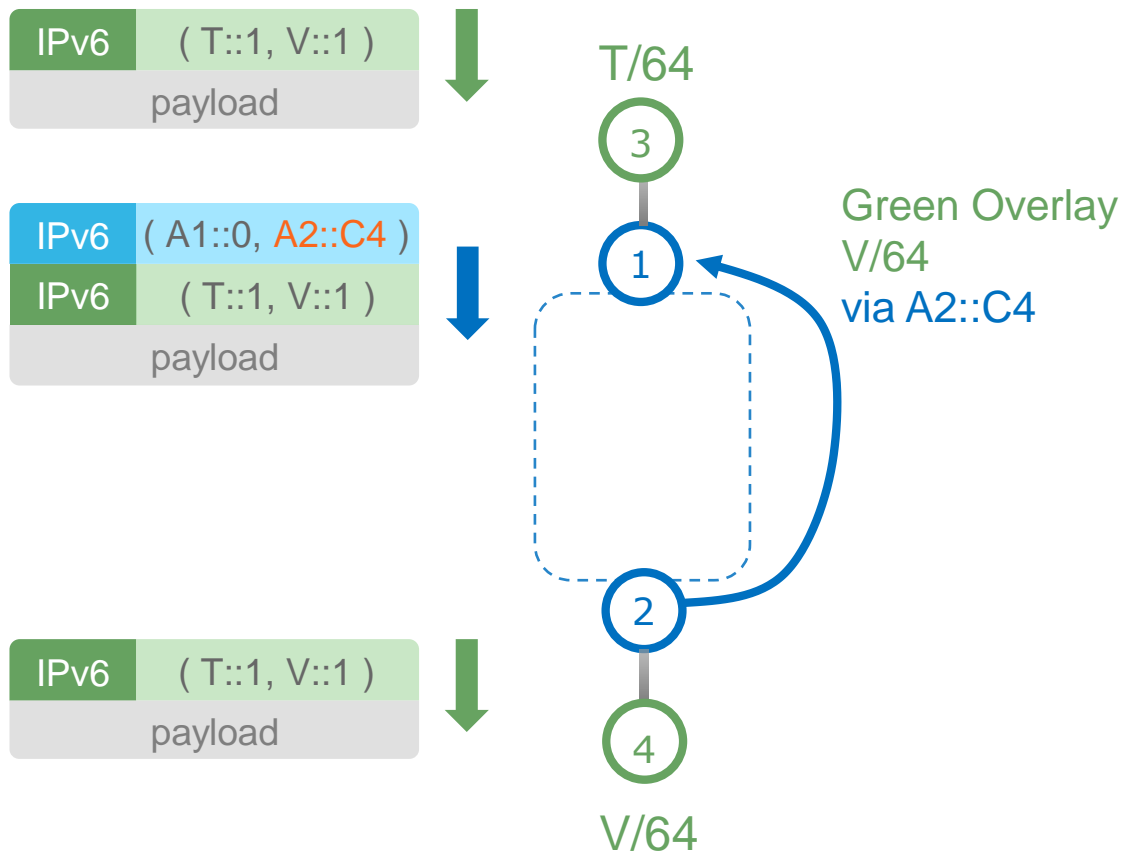
## Algorithm

- SR native



# Overlay

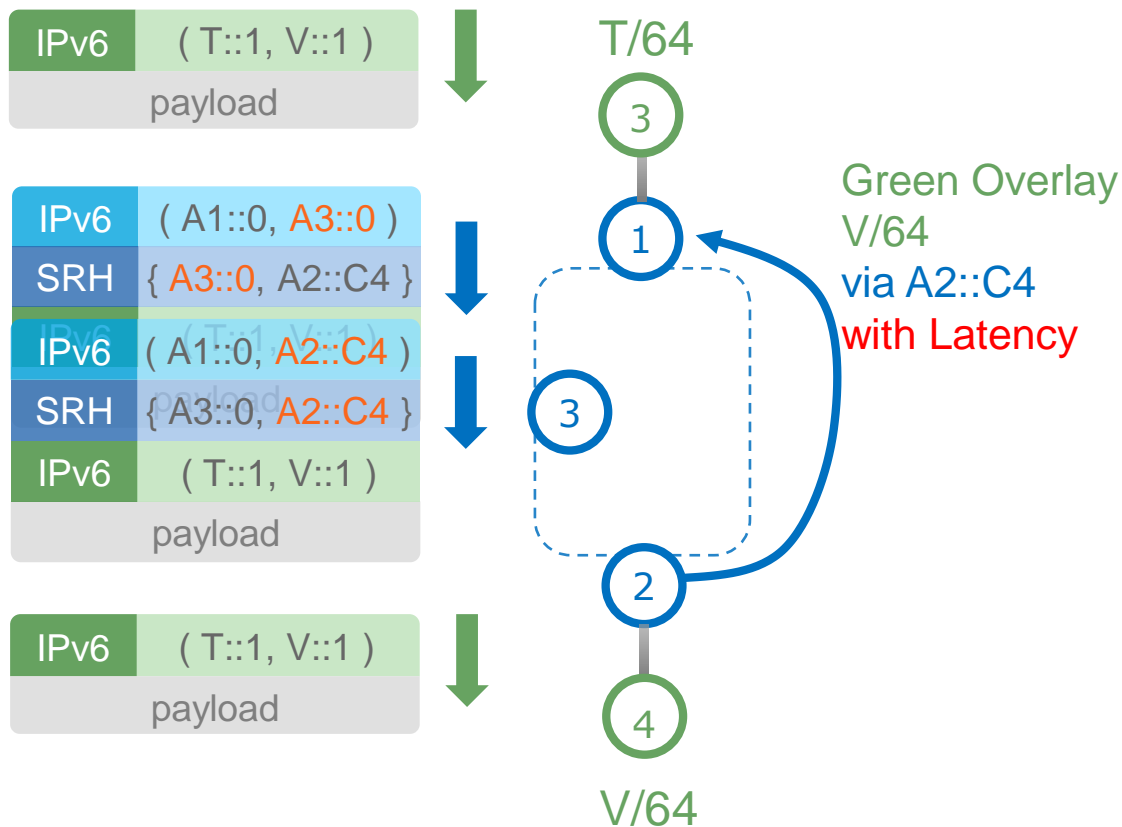
- Automated
  - No tunnel to configure
- Simple
  - Protocol elimination
- Efficient
  - SRv6 for everything



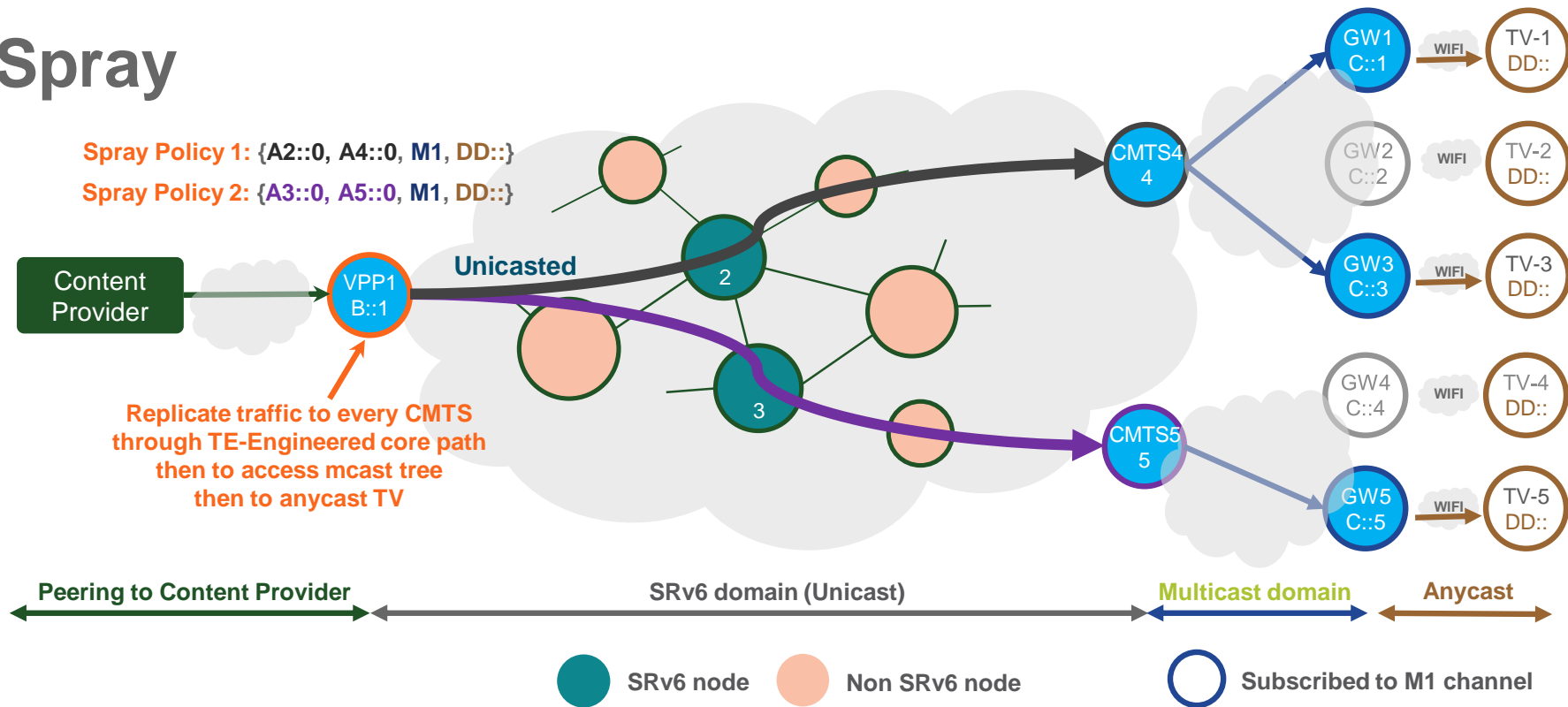


# Overlay with Underlay Control

- SRv6 does not only eliminate unneeded overlay protocols
- SRv6 solves problems that these protocols cannot solve



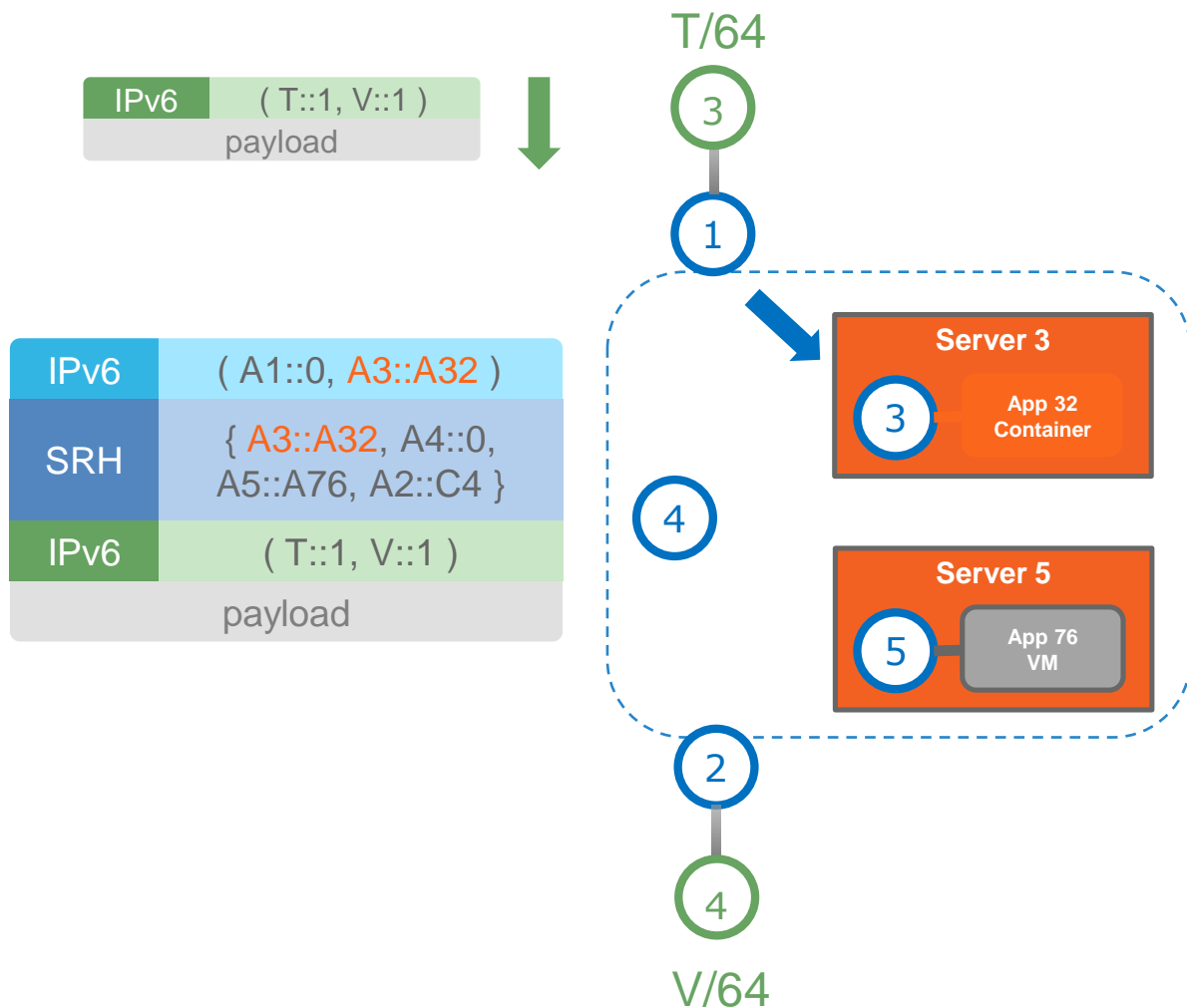
# Spray



Flexible, SLA-enabled and Efficient content injection without multicast core

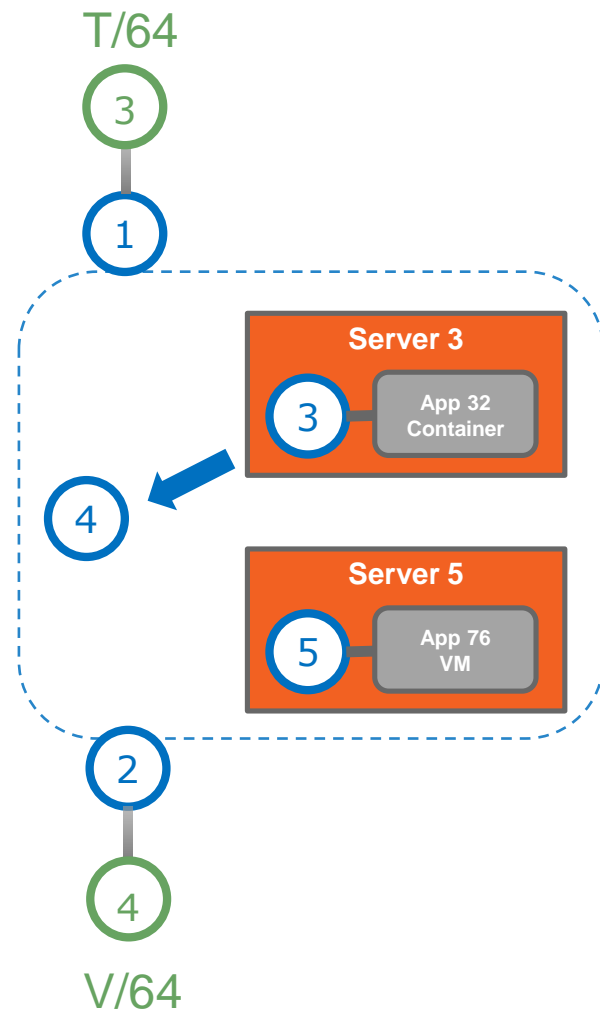
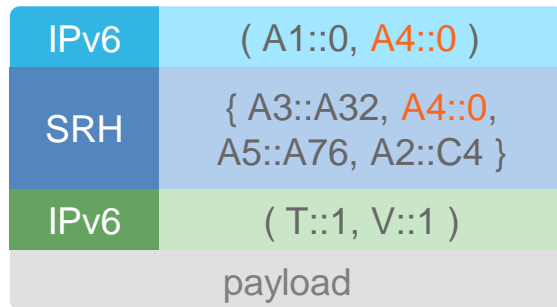
# Integrated NFV

- A3::A32 means
  - App in Container 32
  - @ node A3::/64
- Stateless
  - NSH creates per-chain state in the fabric
  - SR does not
- App is SR aware or not



# Integrated NFV

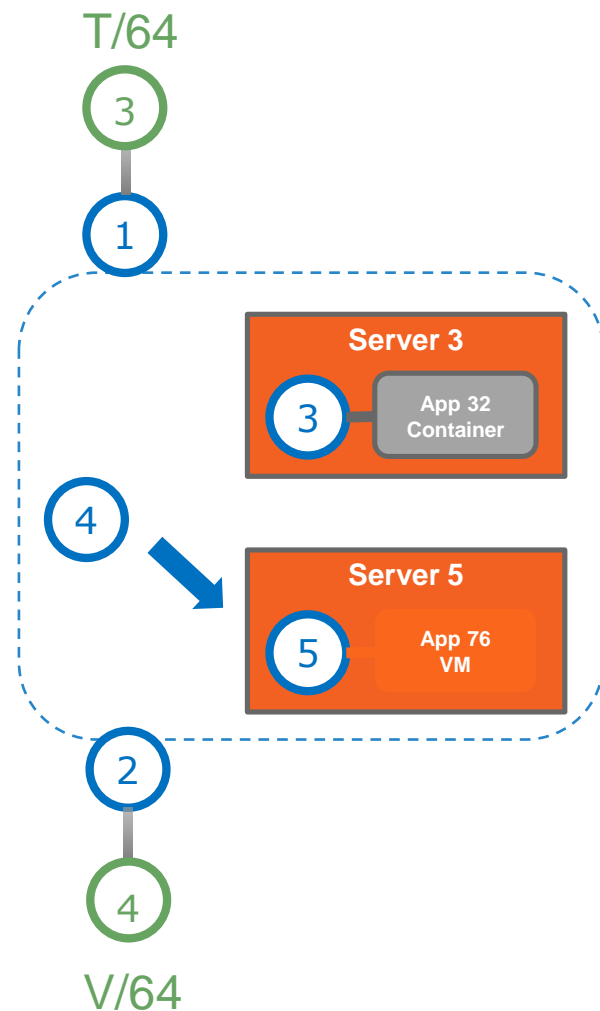
- Integrated with underlay SLA



# Integrated NFV

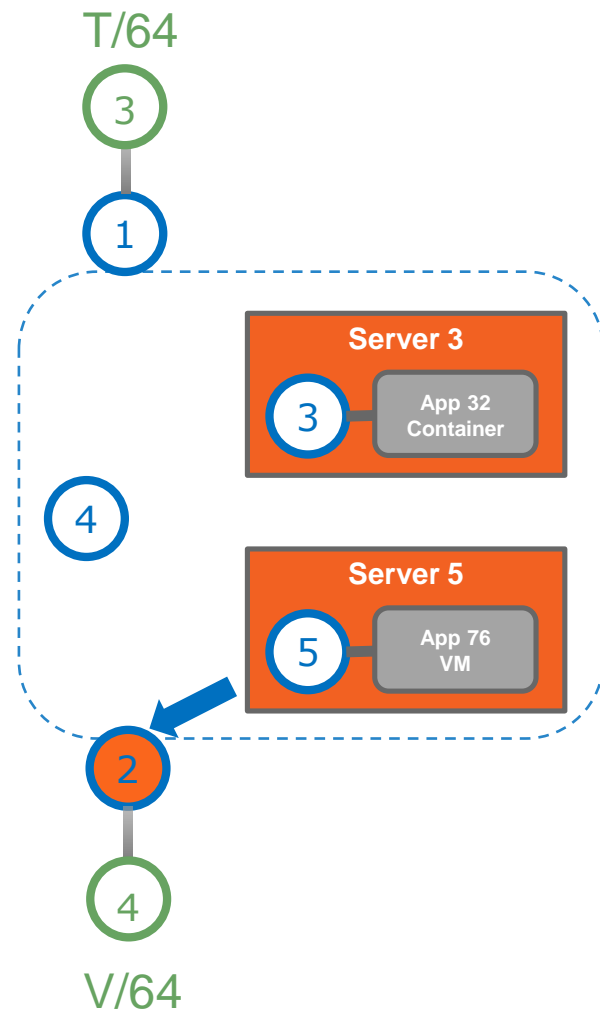
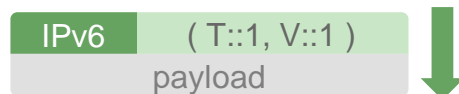
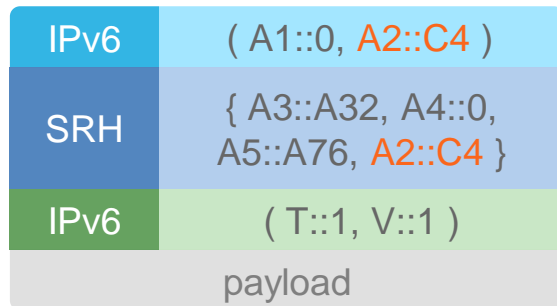
- A5::A76 means
  - App in VM 76
  - @ node A5::/64
- Stateless
  - NSH creates per-chain state in the fabric
  - SR does not
- App is SR aware or not

IPv6	( A1::0, A5::A76 )
SRH	{ A3::A32, A4::0, A5::A76, A2::C4 }
IPv6	( T::1, V::1 )
payload	



# Integrated NFV

- Integrated with Overlay



# More use-cases

- 6CN: enhancing IP to search for Content
- 6LB: enhancing load-balancers
- Video Pipeline
- 5G Slicing
- 5G Ultra-Low Latency

# SRv6 status

- Cisco HW
  - ASR9k - XR
  - ASR1k – XE
- Open-Source
  - Linux 4.10
  - FD.IO





# Conclusion

# Network Programming

- An SRv6 segment is a function at a node
- An SRv6 segment list is a network program
- The network acts as a large computer
- Integrated use-cases well beyond underlay (TE, FRR)
  - NFV
  - Container networking
  - Efficient content management: Spray, 6CN, 6LB
  - Video pipeline
- Simplification: IPv6+SRv6 only !

# SRv6 Leadership

- Bold architecture
- Numerous use-cases
  - FRR, TE, SDN, Overlay with SLA, NFV, Spray, 6CN, 6LB, 5G Slice & LL
- First to demonstrate HW implementation
- First to FCS, field trial and deployment
- Fund university to bring SRv6 in Linux 4.10
- Fund significant SRv6 implementation in FD.IO
- Feel free to join the lead-operator team!

# Complete Your Online Session Evaluation

- Please complete your Online Session Evaluations after each session
- Complete 4 Session Evaluations & the Overall Conference Evaluation (available from Thursday) to receive your Cisco Live T-shirt
- All surveys can be completed via the Cisco Live Mobile App or the Communication Stations



Don't forget: Cisco Live sessions will be available for viewing on-demand after the event at [CiscoLive.com/Online](https://cislive.com/Online)

# Continue Your Education

- Demos in the Cisco campus
- Walk-in Self-Paced Labs
- Lunch & Learn
- Meet the Engineer 1:1 meetings
- Related sessions

Thank You



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Thank You



# Your Time Is Now